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*Technical Specification*

# Digital cellular telecommunications system (Phase 2+);
# General Packet Radio Service (GPRS);
# Overall description of the GPRS radio interface;
# Stage 2
# (3GPP TS 03.64 version 8.12.0 Release 1999)

**GLOBAL SYSTEM FOR**
**MOBILE COMMUNICATIONS**

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# Foreword

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# Contents

# Foreword

This Technical Specification has been produced by the 3$^{rd}$ Generation Partnership Project (3GPP).

The contents of the present document are subject to continuing work within the TSG and may change following formal TSG approval. Should the TSG modify the contents of the present document, it will be re-released by the TSG with an identifying change of release date and an increase in version number as follows:

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# 1 Scope

The present document provides the overall description for lower-layer functions of the General Packet Radio Service (GPRS and EGPRS)) radio interface (Um). ). Within this TS the term GPRS refers to GPRS and EGPRS unless explicitly stated otherwise.

The overall description provides the following information:

- The services offered to higher-layer functions,

- The distribution of required functions into functional groups,

- A definition of the capabilities of each functional group,

- Service primitives for each functional group, including a description of what services and information flows are to be provided, and

- A model of operation for information flows within and between the functions.

The present document is applicable to the following GPRS Um functional layers:

- Radio Link Control functions,

- Medium Access Control functions, and

- Physical Link Control functions.

The present document describes the information transfer and control functions to be used across the radio (Um) interface for communication between the MS and the Network, see Figure 1.

3GPP TS 03.60 [3] describes the overall GPRS logical architecture and the GPRS functional layers above the Radio Link Control and Medium Access Control layer.

3GPP TS 04.07 [5] contains a description in general terms of the structured functions and procedures of this protocol and the relationship of this protocol with other layers and entities.

3GPP TS 04.08 [6] contains the definition of GPRS RLC/MAC procedures when operating on the Common Control Channel (CCCH).

3GPP TS 04.60 [7] contains the definition of RLC/MAC functions when operating on a Packet Data Channel (PDCH).

3GPP TS 04.64 [8] contains functional procedures for the Logical Link Control (LLC) layer above the RLC/MAC.

3GPP TS 05 series defines the Physical Link layer and Physical RF layer.
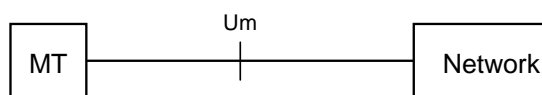


**Figure 1: Scope of GPRS Logical Radio Interface Architecture**

# 2 References

The following documents contain provisions which, through reference in this text, constitute provisions of the present document.

- References are either specific (identified by date of publication, edition number, version number, etc.) or non-specific.

- For a specific reference, subsequent revisions do not apply.

- For a non-specific reference, the latest version applies. In the case of a reference to a 3GPP document (including a GSM document), a non-specific reference implicitly refers to the latest version of that document *in the same Release as the present document.*

[1]          3GPP TR 01.04 (ETR 350): "Digital cellular telecommunications system (Phase 2+); Abbreviations and acronyms".

[2]          3GPP TS 02.60: "Digital cellular telecommunications system (Phase 2+); General Packet Radio Service (GPRS); Stage 2 ".

[3]          3GPP TS 03.60: "Digital cellular telecommunications system (Phase 2+); Stage 2 Service Description of the General Packet Radio Service (GPRS)".

[4]          3GPP TS 04.04: "Digital cellular telecommunications system; Layer 1; General requirements".

[5]          3GPP TS 04.07: "Digital cellular telecommunications system (Phase 2+); Mobile radio interface signalling layer 3 General aspects"

[6]          3GPP TS 04.08: "Digital cellular telecommunications system (Phase 2+); Mobile radio interface layer 3 specification"

[7]          3GPP TS 04.60: "Digital cellular telecommunications system(Phase 2+); General Packet Radio Service (GPRS);  Mobile Station (MS) – Base Station System (BSS) interface; Radio Link Control/Medium Access Control (RLC/MAC) protocol".

[8]          3GPP TS 04.64: "Digital cellular telecommunications system(Phase 2+); General Packet Radio Service (GPRS); Logical Link Control (LLC)".

[9]          3GPP TS 04.65: "Digital cellular telecommunications system (Phase 2+); General Packet Radio Service (GPRS); Subnetwork Dependent Convergence Protocol (SNDCP)".

[10]          3GPP TS 05.01: "Digital cellular telecommunications system (Phase 2+); Physical layer on the radio path, General description".

[11]          3GPP TS 05.02: "Digital cellular telecommunications system (Phase 2+); Multiplexing and multiple access on the radio path".

[12]          3GPP TS 05.03: "Digital cellular telecommunications system (Phase 2+); Channel coding".

[13]          3GPP TS 05.04: "Digital cellular telecommunications system (Phase 2+); Modulation".

[14]          3GPP TS 05.05: "Digital cellular telecommunications system (Phase 2+); Radio transmission and reception".

[15]          3GPP TS 05.08: "Digital cellular telecommunications system (Phase 2+); Radio subsystem link control".

[16]          3GPP TS 05.10: "Digital cellular telecommunications system (Phase 2+); Radio subsystem synchronisation".

# 3     Abbreviations, symbols and definitions

## 3.1     Abbreviations

In addition to abbreviations in 3GPP TR 01.04 [1] and 3GPP TS 02.60 [2] the following abbreviations apply:

| | |
|---|---|
| ARQ | Automatic Repeat reQuest |
| BCS | Block Check Sequence |
| BEC | Backward Error Correction |
| BH | Block Header |
| CFCCH | Compact Frequency Correction Channel |
| CPAGCH | Compact Packet Access Grant Channel |

| | |
|---|---|
| CPBCCH | Compact Packet Broadcast Control Channel |
| CPCCCH | Compact Packet Common Control Channel |
| CPNCH | Compact Packet Notification Channel (for PTM-M on CPCCCH) |
| CPPCH | Compact Packet Paging Channel |
| CPRACH | Compact Packet Random Access Channel |
| CSCH | Compact Synchronization Channel |
| CS | Coding Scheme |
| CU | Cell Update |
| DTM | Dual Transfer Mode |
| EGPRS | Enhanced GPRS |
| FBI | Final Block Indicator |
| FH | Frame Header |
| GGSN | Gateway GPRS Support Node |
| HCS | Header Check Sequence |
| IR | Incremental Redundancy |
| LLC | Logical Link Control |
| MAC | Medium Access Control |
| MCS | Modulation and Coding Scheme |
| NCH | Notification Channel (for PTM-M on CCCH) |
| NSS | Network and Switching Subsystem |
| PACCH | Packet Associate Control Channel |
| PAGCH | Packet Access Grant Channel |
| PBCCH | Packet Broadcast Control Channel |
| PC | Power Control |
| PCCCH | Packet Common Control Channel |
| PDCH | Packet Data Channel |
| PDTCH | Packet Data Traffic Channel |
| PDU | Protocol Data Unit |
| PL | Physical Link |
| PNCH | Packet Notification Channel (for PTM-M on PCCCH) |
| PPCH | Packet Paging Channel |
| PRACH | Packet Random Access Channel |
| PSI | Packet System Information |
| PTCCH | Packet Timing Advance Control Channel |
| RLC | Radio Link Control |
| SGSN | Serving GPRS Support Node |
| SNDC | Subnetwork Dependent Convergence |
| TA | Timing Advance |
| TBF | Temporary Block Flow |
| TFI | Temporary Frame Identity |
| USF | Uplink State Flag |

## 3.2    Symbols

For the purposes of the present document, the following symbols apply:

| | |
|---|---|
| Gb | Interface between an SGSN and a BSC. |
| Um | Interface between MS and GPRS fixed network part. The Um interface is the GPRS network interface for providing packet data services over the radio to the MS. |

## 3.3    Definitions

### 3.3.1    General

GPRS specific definitions can be found in 02.60 [2] and 03.60 [3].

## 3.3.2 EGPRS mobile station

An EGPRS mobile station is a GPRS mobile station with additional capabilities for new radio access protocol features and new modulation and coding schemes. An EGPRS mobile station shall comply with GPRS requirements and the additional requirements defined for an EGPRS mobile station. The support of EGPRS is optional for the mobile station and the network.

## 3.3.3 Dual Transfer Mode

In dual transfer mode, the mobile station is allocated resources providing an RR connection and a Temporary Block Flow on one or more physical channels. This feature is optional for the mobile station and the network. It is only applicable for a mobile station supporting GPRS or EGPRS. Dual transfer mode is a subset of class A mode of operation, which is only possible if there is radio resource allocation co-ordination in the network.

# 4 Packet data logical channels

NOTE: The text in this clause is informative. The normative text is in 3GPP TS 05.02 [11]. Where there is a conflict between these descriptions, the normative text has precedence.

## 4.1 General

This subclause describes the packet data logical channels that are supported by the radio subsystem. The packet data logical channels are mapped onto the physical channels that are dedicated to packet data.

The physical channel dedicated to packet data traffic is called a Packet Data Channel (PDCH).

## 4.2 Packet Common Control Channel (PCCCH) and Compact (CPCCCH)

PCCCH and CPCCCH comprises logical channels for common control signalling used for packet data as described in the following subclauses.

### 4.2.1 Packet Random Access Channel (PRACH) and Compact Packet Random Access Channel (CPRACH) - uplink only

PRACH and CPRACH are used by MS to initiate uplink transfer for sending data or signalling information. Packet Access burst and Extended Packet Access burst are used on PRACH. Extended Packet Access burst is used on CPRACH.

### 4.2.2 Packet Paging Channel (PPCH) and Compact Packet Paging Channel (CPPCH) - downlink only

PPCH and CPPCH are used to page an MS prior to downlink packet transfer. PPCH and CPPCH use paging groups in order to allow usage of DRX mode. PPCH can be used for paging of both circuit switched and packet data services. The paging for circuit switched services on PPCH is applicable for class A and B GPRS MSs in Network operation mode I, see 3GPP TS 03.60 [3].

### 4.2.3 Packet Access Grant Channel (PAGCH) and Compact Packet Access Grant Channel (CPAGCH) - downlink only

PAGCH and CPAGCH are used in the packet transfer establishment phase to send resource assignment to an MS prior to packet transfer.

### 4.2.4 Packet Notification Channel (PNCH) and Compact Packet Notification Channel (CPNCH) - downlink only

PNCH and CPNCH are used to send a PTM-M (Point To Multipoint - Multicast) notification to a group of MSs prior to a PTM-M packet transfer.

DRX mode shall be provided for monitoring PNCH and CPNCH. Furthermore, a "PTM-M new message" indicator may optionally be sent on all individual paging channels to inform MSs interested in PTM-M when they need to listen to PNCH and CPNCH.

The PTM-M service is not specified in GPRS Phase 1.

## 4.3 Packet Broadcast Control Channel (PBCCH) and Compact Packet Broadcast Control Channel (CPBCCH) - downlink only

PBCCH and CPBCCH broadcast packet data specific System Information. If PBCCH is not allocated, the packet data specific system information is broadcast on BCCH. For Compact, CPBCCH shall be allocated. CPBCCH and  BCCH are mutually exclusive.

## 4.4 Packet Traffic Channels

### 4.4.1 Packet Data Traffic Channel (PDTCH)

PDTCH is a channel allocated for data transfer. It is temporarily dedicated to one MS or to a group of MSs in the PTM-M case. In the multislot operation, one MS may use multiple PDTCHs in parallel for individual packet transfer.

All packet data traffic channels are uni-directional, either uplink (PDTCH/U), for a mobile originated packet transfer or downlink (PDTCH/D) for a mobile terminated packet transfer.

## 4.5 Packet Dedicated Control Channels

### 4.5.1 Packet Associated Control Channel (PACCH)

PACCH conveys signalling information related to a given MS. The signalling information includes e.g. acknowledgements and power control information. PACCH carries also resource assignment and reassignment messages, comprising the assignment of a capacity for PDTCH(s) and for further occurrences of PACCH. The PACCH shares resources with PDTCHs,  that are currently assigned to one MS. Additionally, an MS that is currently involved in packet transfer, can be paged for circuit switched services on PACCH.

### 4.5.2 Packet Timing advance Control Channel, uplink (PTCCH/U)

PTCCH/U is used to transmit random access burst to allow estimation of the timing advance for one MS in packet transfer mode.

PTCCH/U shall not be used for DTM.

### 4.5.3 Packet Timing advance Control Channel, downlink (PTCCH/D)

PTCCH/D is used to transmit timing advance information updates to several MS. One PTCCH/D is paired with several PTCCH/U's.

PTCCH/D shall be ignored by MS operating in DTM.

# 5 Mapping of packet data logical channels onto physical channels

NOTE: The text in this clause is informative. The normative text is in 3GPP TS 05.02 [11]. Where there is a conflict between these descriptions, the normative text has precedence.

## 5.1 General

Different packet data logical channels can occur on the same physical channel (i.e. PDCH). The sharing of the physical channel is based on blocks of 4 consecutive bursts, except for PTCCH. The mapping in frequency of PDCH on to the physical channel shall be as defined in GSM 05.02 [11].

A PDCH may be either full-rate (PDCH/F) or half-rate (PDCH/H). PDCH/H is only applicable to DTM. See GSM 05.02 [11].

GPRS and EGPRS employ the same physical layer, except for the PDTCH.

On PRACH, CPRACH and PTCCH/U, access bursts are used. On all other packet data logical channels, radio blocks comprising 4 normal bursts are used. The only exception is some messages on uplink PACCH which comprise 4 consecutive access bursts (to increase robustness).

## 5.2 Packet Common Control Channels (PCCCH and CPCCCH)

At a given time, the logical channels of the PCCCH are mapped on different physical resources than the logical channels of the CCCH.

The PCCCH and CPCCCH do not have to be allocated permanently in the cell. Whenever the PCCCH is not allocated, the CCCH shall be used to initiate a packet transfer. For Compact, CPCCCH shall be allocated.

One given MS may use only a subset of the PCCCH and CPCCCH, the subset being mapped onto one physical channel (i.e. PDCH).

The PCCCH, when it exists:

- is mapped on one or several physical channels according to a 52-multiframe, In that case the PCCCH, PBCCH and PDTCH share same physical channels (PDCHs).

The existence and location of the PCCCH shall be broadcast on the cell.

Since GSM phase 1 and phase 2 MS can only see and use the CCCH, the use on the PCCCH can be optimised for GPRS e.g. a PRACH of 11 bits can be used on uplink.

For Compact, one radio frequency channel of the cell allocation shall be used to carry synchronization information and the CPBCCH, this shall be known as the primary Compact carrier. All other radio frequency channels of the cell allocation shall be known as secondary Compact carriers.

For primary and secondary Compact carriers, CPCCCHs shall be allocated on only one timeslot (which is associated with a time group as defined in GSM 05.02 [11]). This time group is known as the serving time group and rotates over odd timeslot numbers as follows: 7, 5, 3, 1, 7, 5, …. The CPCCCH is mapped according to a Compact 52-multiframe and the serving time group rotation occurs between frame numbers (FN) mod 52 = 3 and 4.

### 5.2.1 Packet Random Access Channel (PRACH and CPRACH)

The PRACHand CPRACH are mapped on one or several physical channels. The physical channels on which the PRACH is mapped are derived by the MS from information broadcast on the PBCCH or BCCH. The physical channels on which the CPRACH is mapped are derived by the MS from information broadcast on the CPBCCH.

PRACH and CPRACH are determined by the Uplink State Flag marked as free that is broadcast continuously on the corresponding downlink (see subclause 6.6.4.1). Additionally, a predefined fixed part of the multiframe structure for PDCH can be used as PRACH or CPRACH only and the information about the mapping on the physical channel is

broadcast on PBCCH or CPBCCH. During those time periods an MS does not have to monitor the USF that is simultaneously broadcast on the downlink.

## 5.2.2     Packet Paging Channel (PPCH and CPPCH)

The PPCH and CPPCH are mapped on one or several physical channels. The exact mapping on each physical channel follows a predefined rule (see subclause 6.1.2), as it is done for the PCH.

The physical channels on which the PPCH or CPPCH are mapped, as well as the rule that is followed on the physical channels, are derived by the MS from information broadcast on the PBCCH or CPBCCH.

## 5.2.3     Packet Access Grant Channel (PAGCH and CPAGCH)

The PAGCH and CPAGCH are mapped on one or several physical channels. The exact mapping on each physical channel follows a predefined rule (see subclause 6.1.2).

The physical channels on which the PAGCH or CPAGCH are mapped, as well as the rule that is followed on the physical channels, are derived by the MS from information broadcast on the PBCCH or CPBCCH.

## 5.2.4     Packet Notification Channel (PNCH and CPNCH)

The PNCH and CPNCH are mapped on one or several blocks on PCCCH and CPCCCH. The exact mapping follows a predefined rule. The mapping is derived by the MS from information broadcast on the PBCCH or CPBCCH.

## 5.3     Packet Broadcast Control Channel (PBCCH and CPBCCH)

The PBCCH and CPBCCH shall be mapped on one or several physical channels. The exact mapping on each physical channel follows a predefined rule (see subclause 6.1.2), as it is done for the BCCH. For Compact, CPBCCH shall be allocated. CPBCCH and BCCH are mutually exclusive.

The existence of the PCCCH, and consequently the existence of the PBCCH, is indicated on the BCCH.

For Compact, one radio frequency channel of the cell allocation shall be used to carry synchronization information and the CPBCCH, this shall be known as the primary Compact carrier. All other radio frequency channels of the cell allocation shall be known as secondary Compact carriers.

The CPBCCH shall be mapped on only one timeslot (which is associated with a time group as defined in GSM 05.02 [11]). This time group is known as the serving time group and rotates over odd timeslot numbers as follows: 7, 5, 3, 1, 7, 5, …. The CPBCCH is mapped according to a Compact 52-multiframe and the serving time group rotation occurs between frame numbers (FN) mod 52 = 3 and 4. The exact mapping follows a predefined rule (see subclause 6.1.2).

## 5.3a     Compact Frequency Correction Channel (CFCCH)

The CFCCH is the same as the FCCH with one exception — the FCCH is mapped onto a 51-multiframe as defined in GSM 05.02 [11].

## 5.3b     Compact Synchronization Channel (CSCH)

The CSCH is similar to the SCH. The major difference is that the SCH is mapped onto a 51-multiframe as defined in GSM 05.02 [11]. This results in a different layout for the reduced TDMA frame number (RFN).

## 5.4     Packet Timing advance Control Channel (PTCCH)

Two defined frames of multiframe are used to carry PTCCH (see subclause 6.1.2). The exact mapping of PTCCH/U sub-channels and PTCCH/D shall be as defined in GSM 05.02 [11].

On PTCCH/U, access bursts are used. On PTCCH/D, four normal bursts comprising a radio block are used.

# 5.5 Packet Traffic Channels

## 5.5.1 Packet Data Traffic Channel (PDTCH)

One PDTCH is mapped onto one physical channel.

Up to eight PDTCHs, with different timeslots but with the same frequency parameters, may be allocated to one MS at the same time.

## 5.5.2 Packet Associated Control Channel (PACCH)

PACCH is dynamically allocated on the block basis on the same physical channel as carrying PDTCHs.However, one block PACCH allocation is used on the physical channel carrying only PCCCH, when the MS is polled to acknowledge the initial assignment message.

PACCH is of a bi-directional nature, i.e. it can dynamically be allocated both on the uplink and on the downlink regardless on whether the corresponding PDTCH assignment is for uplink or downlink.

When PDTCH(s) is assigned on the uplink, the corresponding downlink timeslots have continuously to be monitored by the MS for possible occurrences of PACCH. The MS can use the uplink assignment for sending PACCH blocks whenever needed. In case of extended dynamic allocation (see subclause 6.6.4.4), if the resource assigned by the network does not allow the multislot MS (see GSM 05.02 [11] , annex B) to monitor the USF on all the assigned PDCHs, the PACCH blocks shall be mapped on one PDCH in the list of assigned PDCHs.

When PDTCH(s) is assigned on the downlink, every occurrence of an uplink PACCH block is determined by polling in one of the preceding downlink blocks (transferred on the same PDCH). The network can use the downlink assignment for sending PACCH blocks whenever needed.

During an uplink allocation a MS using a fixed allocation (see subclause 6.6.4.4) must monitor the assigned PACCH timeslot during all blocks where the uplink is unassigned a number ofconsecutive timeslots. The number of consecutive timeslots depends upon the multislot class of the MS. The network shall transmit a PACCH block to a MS using a fixed allocation only during the same size timeslot gap in the uplink allocation on the PACCH.

During a downlink transmission the network shall not send downlink data to a MS during uplink PACCH timeslots or in a number of timeslot preceding and following the uplink PACCH block. The number of timeslot preceding and following the uplink PACCH timeslots depends upon the multislot class of the half duplex MS.

# 5.6 Downlink resource sharing

Different packet data logical channels can be multiplexed on the downlink on the same physical channel (i.e. PDCH). See details in GSM 05.02 [11]. The type of message which is indicated in the radio block header allows differentiation between the logical channels. Additionally, the MS identity allows differentiation between PDTCHs and PACCHs assigned to different MSs.

In addition, in dual transfer mode the network may allocate a PDCH dedicated to the MS. Even in the case of exclusive allocation, the network shall use the MS identity and the type of message in the radio block header.

# 5.7 Uplink resource sharing

Different packet data logical channels can be multiplexed on the uplink of the same physical channel (i.e. PDCH). See details in GSM 05.02 [11]. The type of message which is indicated in the radio block header, allows differentiation between the logical channels. Additionally, the MS identity allows differentiation between PDTCHs and PACCHs assigned to different MSs.

In addition, in dual transfer mode the network may allocate a PDCH dedicated to the MS.

# 6 Radio Interface (Um)

The logical architecture of the GPRS Um interface can be described using a reference model consisting of functional layers as shown in Figure 3. Layering provides a mechanism for partitioning communications functions into manageable subsets.

Communication between the MS and the Network occurs at the Physical RF, Physical Link, Radio Link Control/Medium Access Control (RLC/MAC), Logical Link Control (LLC) and Subnetwork Dependent Convergence layers.

## 6.1 Radio Resource management principles

### 6.1.1 Allocation of resources for the GPRS

A cell supporting GPRS may allocate resources on one or several physical channels in order to support the GPRS traffic. Those physical channels (i.e. PDCHs), shared by the GPRS MSs, are taken from the common pool of physical channels available in the cell. The allocation of physical channels to circuit switched services and GPRS is done dynamically according to the "capacity on demand" principles described below.

Common control signalling required by GPRS in the initial phase of the packet transfer is conveyed on PCCCH, when allocated, or on CCCH. This allows the operator to have capacity allocated specifically to GPRS in the cell only when a packet is to be transferred.

For Compact, common control signaling required by the mobile station in the initial phase of the packet transfer is conveyed on CPCCCH.

#### 6.1.1.1 Master-Slave concept

At least one PDCH, acting as a master, accommodates packet common control channels that carry all the necessary control signalling for initiating packet transfer (i.e. PCCCH), whenever that signalling is not carried by the existing CCCH, as well as user data and dedicated signalling (i.e. PDTCH and PACCH). Other PDCHs, acting as slaves, are used for user data transfer and for dedicated signalling.

For Compact, one radio frequency channel of the cell allocation shall be used to carry synchronization information and the CPBCCH, this shall be known as the primary Compact carrier. All other radio frequency channels of the cell allocation shall be known as secondary Compact carriers.

For the primary Compact carrier, timeslot numbers (TN) 1, 3, 5, and 7, acting as a master, accommodate packet common control channels that carry all necessary control signalling for initiating packet transfer as well as user data and dedicated signalling (i.e., PDTCH and PACCH). TNs 0, 2, 4, and 6, acting as slaves, are used for user data transfer and for dedicated signalling.

For the secondary Compact carrier(s) carrying CPCCCH, timeslot numbers (TN) 1, 3, 5, and 7, acting as a master, accommodate packet common control channels that carry all necessary control signalling for initiating packet transfer as well as user data and dedicated signalling. TNs 0, 2, 4, and 6, acting as slaves, are used for user data transfer and for dedicated signalling.

For the secondary Compact carrier(s) not carrying CPCCCH, timeslot numbers (TN) 0 through 7, acting as slaves, are used for user data transfer and for dedicated signalling.

#### 6.1.1.2 Capacity on demand concept

The GPRS does not require permanently allocated PDCHs. The allocation of capacity for GPRS can be based on the needs for actual packet transfers which is here referred to as the "capacity on demand" principle. The operator can, as well, decide to dedicate permanently or temporarily some physical resources (i.e. PDCHs) for the GPRS traffic.

When the PDCHs are congested due to the GPRS traffic load and more resources are available in the cell, the Network can allocate more physical channels as PDCHs.

However, the existence of PDCH(s) does not imply the existence of PCCCH.

When no PCCCH is allocated in a cell, all GPRS attached MSs camp on the CCCH.

In response to a Packet Channel Request sent on CCCH from the MS that wants to transmit GPRS packets, the network can assign resources on PDCH(s) for the uplink transfer.. After the transfer, the MS returns to CCCH.

When PCCCH is allocated in a cell, all GPRS attached MSs camp on it. PCCCH can be allocated either as the result of the increased demand for packet data transfers or whenever there is enough available physical channels in a cell (to increase the quality of service). The information about PCCCH is broadcast on BCCH. When the PCCCH capacity is inadequate, it is possible to allocate additional PCCCH resources on one or several PDCHs. If the network releases the last PCCCH, the MS performs cell re-selection.

For Compact, CPBCCH shall be allocated. CPBCCH is a stand-alone packet control channel for Compact. CPCCCH shall be allocated. The information about CPCCCH is broadcast on CPBCCH. When CPCCCH capacity is inadequate, it is possible to allocate additional CPCCCH resources on primary and secondary Compact carriers.

### 6.1.1.3 Procedures to support capacity on demand

The number of allocated PDCHs in a cell can be increased or decreased according to demand. The following principles can be used for the allocation:

- Load supervision:

  A load supervision function may monitor the load of the PDCHs and the number of allocated PDCHs in a cell can be increased or decreased according to demand. Load supervision function may be implemented as a part of the Medium Access Control (MAC) functionality. The common channel allocation function located in BSC is used for the GSM services.

- Dynamic allocation of PDCHs:

  Unused channels can be allocated as PDCHs to increase the overall quality of service for GPRS.

  Upon resource demand for other services with higher priority, de-allocation of PDCHs can take place.

### 6.1.1.4 Release of PDCH not carrying PCCCH

The fast release of PDCH is an important feature for possibility to dynamically share the same pool of radio resources for packet and circuit-switched services.

There are following possibilities:

- **Wait for all the assignments to terminate on that PDCH**

- **Individually notify all the users that have assignment on that PDCH**

  Packet Uplink Assignment and Packet Downlink Assignment messages can be used for that purpose. The network side has to send such notifications on PACCH(s) individually to each affected MS.

- **Broadcast the notification about de-allocation**

  Simple and fast method to broadcast the Packet PDCH Release on all the PDCHs lying on the same carrier as the PDCH to be released. All MSs monitor the possible occurrences of PACCH on one channel and should capture such notification.

In practice, a combination of all the methods can be used.

There may occur the case where an MS remains unaware of the released PDCH. In that case, such MS may cause some interference when wrongly assuming that the decoded Uplink State Flag (see Subclause 6.6.4.1.) denotes the following uplink block period reserved to it. After not getting proper response from the network, the MS would self break the RLC connection.

## 6.1.2 Multiframe structure for PDCH

NOTE: The text in this clause is informative. The normative text is in GSM 05.02 [11]. Where there is a conflict between these descriptions, the normative text has precedence.

The mapping in time of the logical channels is defined by a multiframe structure. The multiframe structure for PDCH consists of 52 TDMA frames, divided into 12 blocks (of 4 frames), 2 idle frames and 2 frames used for the PTCCH according to Figure 2.

52 TDMA Frames

| B0 | B1 | B2 | T | B3 | B4 | B5 | X | B6 | B7 | B8 | T | B9 | B10 | B11 | X |

X = Idle frame
T = Frame used for PTCCH
B0 - B11 = Radio blocks
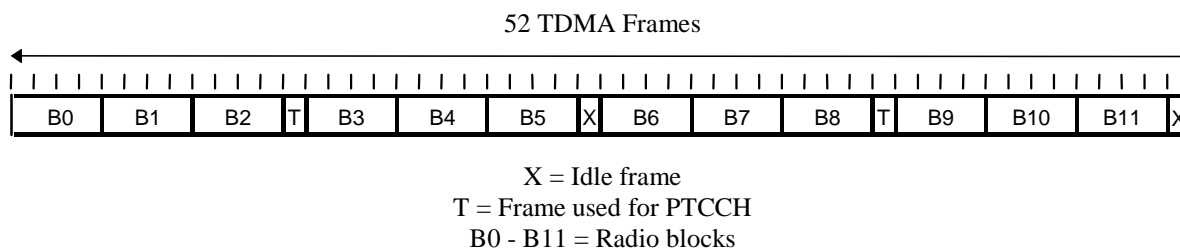
**Figure 2: Multiframe structure for PDCH**

The mapping of logical channels onto the radio blocks is defined in the rest of this subclause by means of the ordered list of blocks (B0, B6, B3, B9, B1, B7, B4, B10, B2, B8, B5, B11).

One PDCH that contains PCCCH (if any) is indicated on BCCH. That PDCH is the only one that contains PBCCH blocks. On the downlink of this PDCH, the first block (B0) in the ordered list of blocks is used as PBCCH. If required, up to 3 more blocks on the same PDCH can be used as additional PBCCH. Any additional PDCH containing PCCCH is indicated on PBCCH.

On any PDCH with PCCCH (with or without PBCCH), the next up to 12 blocks in the ordered list of blocks are used for PAGCH, PNCH, PDTCH or PACCH in the downlink. The remaining blocks in the ordered list are used for PPCH, PAGCH, PNCH, PDTCH or PACCH in the downlink. In all cases, the actual usage of the blocks is indicated by the message type. On an uplink PDCH that contains PCCCH, all blocks in the multiframe can be used as PRACH, PDTCH or PACCH. Optionally, the first blocks in the ordered list of blocks can only used as PRACH. The MS may chose to either ignore the USF (consider it as FREE) or use the USF to determine the PRACH in the same way as for the other blocks.

The mapping of channels on multiframes are controlled by several parameters broadcast on PBCCH.

On a PDCH that does not contain PCCCH, all blocks can be used as PDTCH or PACCH. The actual usage is indicated by the message type.

Two frames are used for PTCCH (see GSM 05.02 [11] ) and the two idle frames as well as the PTCCH frames can be used by the MS for signal measurements and BSIC identification.

## 6.1.2a  Multiframe structure for Compact PDCH

NOTE:   The text in this clause is informative. The normative text is in GSM 05.02 [11]. Where there is a conflict between these descriptions, the normative text has precedence.

For Compact, one radio frequency channel of the cell allocation shall be used to carry synchronization information and the CPBCCH, this shall be known as the primary Compact carrier. All other radio frequency channels of the cell allocation shall be known as secondary Compact carriers.

For the primary Compact carrier, timeslot numbers (TN) 1, 3, 5, and 7 accommodate packet common control channels (i.e., CPBCCH and CPCCCH) as well as user data and dedicated signalling (i.e., PDTCH and PACCH). TNs 0, 2, 4, and 6 are used for user data transfer and for dedicated signalling.

For the secondary Compact carrier(s) carrying CPCCCH, timeslot numbers (TN) 1, 3, 5, and 7 accommodate packet common control channels as well as user data and dedicated signalling. TNs 0, 2, 4, and 6 are used for user data transfer and for dedicated signalling.

For the secondary Compact carrier(s) not carrying CPCCCH, timeslot numbers (TN) 0 through 7 are used for user data transfer and for dedicated signalling.

For Compact, a base station is typically assigned at least 3 frequencies (one per cell which translates into one primary Compact carrier per cell allocation) using a 1/3 frequency re-use pattern. Each cell is assigned one time group based upon which timeslot number is allocated for control (see GSM 05.02 [11]). This is known as the serving time group.

Timeslot mapping and rotation of the control channels is used such that control channels belonging to a serving time group are rotated over odd timeslot numbers as follows: 7, 5, 3, 1, 7, 5 … . The rotation occurs between frame numbers (FN) mod 52 = 3 and 4. Packet switched logical channels PDTCH, PACCH, and PTCCH are never rotated.

For Compact, packet switched logical channels are mapped onto a Compact 52-multiframe. A Compact 52-multiframe consists of 12 blocks of 4 consecutive frames, 2 idle frames (which can be used for CFCCH and CSCH), and 2 frames used for PTCCH (see GSM 05.02 [11] and 05.10 [16]) as shown in Figure 2. A block allocated to a given logical channel comprises one radio block or, in uplink only, 4 random access bursts. The type of channel may vary on a block by block basis.

The mapping of CPBCCH onto the radio blocks is defined by means of the ordered list of blocks (B0, B6, B3, B9, B1, B7, B4, B10, B2, B8, B5, B11). On the downlink of the primary Compact carrier, the first block (B0) shall be used as CPBCCH. If required, up to 3 more blocks on the primary Compact carrier can be used as additional CPBCCH. The next up to 12 blocks in the ordered list of blocks are used for CPAGCH, CPNCH, PDTCH, and PACCH in the downlink. The remaining blocks in the ordered list are used for CPPCH, CPAGCH, and CPNCH in the downlink. In all cases, the actual usage of the blocks is indicated by the message type. The same applies to secondary Compact carriers.

In the uplink of the primary Compact carrier and secondary Compact carrier(s), all blocks in the multiframe can be used as CPRACH. However, a prioritization scheme is recommended (see GSM 05.02 [11]). The MS may chose to either ignore the USF (consider it as FREE) or use the USF to determine the CPRACH in the same way as for the other blocks. Optionally, the first blocks in the ordered list of blocks can only be used as CPRACH.

The mapping of channels on multiframes are controlled by several parameters broadcast on CPBCCH.

## 6.1.2b    Multiframe structure for PDCH/H

NOTE:    The text in this clause is informative. The normative text is in GSM 05.02 [11]. Where there is a conflict between these descriptions, the normative text has precedence.

The mapping in time of the logical channels is defined by a multiframe structure. The multiframe structure for PDCH/H consists of 52 TDMA frames, divided into 6 blocks (of 4 frames) and 2 idle frames according to figure 3.. No frames are used for PTCCH (see GSM 05.02 [11]) and the two idle frames can be used by the MS for signal measurements and BSIC identification.
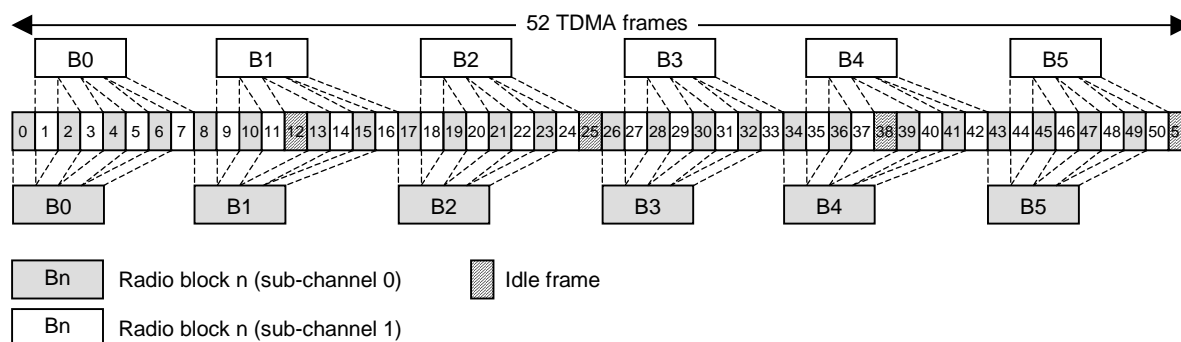


**Figure 3: Multiframe structure for PDCH/H**

A PDCH/H cannot be used as a PCCCH. On a PDCH/H all blocks can be used as PDTCH or PACCH. The actual usage is indicated by the message type. The PDCH/H shall only be allocated to a mobile station in conjunction with a TCH/H in the other subchannel of the physical channel.

## 6.1.3    Scheduling of PBCCH information.

An MS attached to GPRS shall not be required to monitor BCCH if a PBCCH exists. All system information relevant for GPRS and some information relevant for circuit switched services (e.g. the access classes) shall in this case be broadcast on PBCCH. For Compact, CPBCCH shall be allocated. CPBCCH and BCCH are mutually exclusive.

In order to facilitate the MS operation, the network is required to transmit certain types of Packet System Information (PSI) messages in specific multiframes and specific PBCCH or CPBCCH blocks within the multiframes. The exact scheduling is in GSM 05.02 [11].

When no PCCCH is allocated, the MS camps on CCCH and receives all system information on BCCH. Any necessary GPRS specific system information shall in that case be broadcast on BCCH. For Compact, CPCCCH shall be allocated.

## 6.1.4 SMS cell broadcast

The MS reading of the primary and extended CBCH is occasionally interrupted by MS idle mode procedures when the MS is GPRS attached and in packet idle mode.

# 6.2 Radio Resource operating modes

Radio Resource (RR) management procedures are characterised by two different RR operating modes. Each mode describes a certain amount of functionality and information allocated. RR procedures and RR operating modes are specified in GSM 04.07 [5].

## 6.2.1 Packet idle mode

Packet idle mode is not applicable to an MS supporting DTM that has an ongoing RR connection. An MS that supports DTM, that has an ongoing RR connection and that has no allocated packet resource is in dedicated mode.

In packet idle mode no Temporary Block Flow (see subclause 6.6.4.2) exists. Upper layers can require the transfer of a LLC PDU which, implicitly, may trigger the establishment of TBF and transition to packet transfer mode.

In packet idle mode, the MS listens to the PBCCH and to the paging sub-channel for the paging group the MS belongs to in idle mode. If PCCCH is not present in the cell, the mobile station listens to the BCCH and to the relevant paging sub-channels.

While operating in packet idle mode, a mobile station belonging to GPRS MS class A may simultaneously enter the different RR service modes defined in GSM 04.08 [6]. A mobile station belonging to either of GPRS MS class B or C leaves both packet idle mode and packet transfer modes before entering dedicated mode, group receive mode or group transmit mode.

## 6.2.2 Packet transfer mode

Packet transfer mode is not applicable to a mobile station supporting DTM that has an ongoing RR connection. A DTM mobile station with an ongoing RR connection and with packet resources allocated is in dual transfer mode (see 6.2.3).

In packet transfer mode, the mobile station is allocated radio resource providing a Temporary Block Flow on one or more physical channels. Continuous transfer of one or more LLC PDUs is possible. Concurrent TBFs may be established in opposite directions. Transfer of LLC PDUs in RLC acknowledged or RLC unacknowledged mode is provided.

When selecting a new cell, mobile station leaves the packet transfer mode, enters the packet idle mode where it switches to the new cell, read the system information and may then resume to packet transfer mode in the new cell.

While operating in packet transfer mode, a mobile station belonging to GPRS MS class A may simultaneously enter the different RR service modes defined in GSM 04.18. A mobile station belonging to either of GPRS MS class B or C leaves both packet idle mode and packet transfer modes before entering dedicated mode, group receive mode or group transmit mode.

## 6.2.3 Dual transfer mode

In dual transfer mode, the MS has an ongoing RR connection and is allocated radio resource providing a Temporary Block Flow on one or more physical channels. Continuous transfer of one or more LLC PDUs is possible. Concurrent TBFs may be established in opposite directions. Transfer of LLC PDUs in RLC acknowledged or RLC unacknowledged mode is provided.

While in dual transfer mode the MS performs all the tasks of dedicated mode. In addition, upper layers can require:

- the release of all the packet resources, which triggers the transition to dedicated mode.

- the release of the RR resources, which triggers the transition to idle mode and packet idle mode.

When handed over to a new cell, the MS leaves the dual transfer mode, enters the dedicated mode where it switches to the new cell, may read the system information messages sent on the SACCH and may then enter dual transfer mode in the new cell.

## 6.2.4 Correspondence between Radio Resource operating modes and Mobility Management States

The Mobility Management states are defined in GSM 03.60 [3]. Table 1 provides the correspondence between Radio Resource states and Mobility Management states:

**Table 1a: Correspondence between RR operating modes and MM states (non-DTM capable MS)**

| RR BSS | Packet transfer mode | Measurement report reception | No state | No state |
|---|---|---|---|---|
| RR MS | Packet transfer mode | Packet idle mode | | Packet idle mode |
| MM (NSS and MS) | Ready | | | Standby |

**Table 1b: Correspondence between RR operating modes and MM states (DTM capable MS)**

| RR BSS | Dual transfer mode | Dedicated mode | Packet transfer mode | Measurement report reception | No state | Dedicated mode | No state |
|---|---|---|---|---|---|---|---|
| RR MS | Dual transfer mode | Dedicated mode | Packet transfer mode | CS idle and packet idle | | Dedicated mode | CS idle and packet idle |
| GMM (NSS and MS) | Ready | | | | | Standby | |

Each state is protected by a timer. The timers run in the MS and the network.

Packet transfer mode is guarded by RLC protocol timers.

## 6.2.5 Transitions between RR operating modes

The RR modes, and therefore the transitions between them, are different for each mode of operation (see 23.060).

Figure 3a shows the four RR states for an MS in mode of operation A that does not support DTM. The four states can be regarded as the combination of two state machines with two RR states each:

- on the circuit switched part, idle mode and dedicated mode

- on the GPRS part, packet idle mode and packet transfer mode
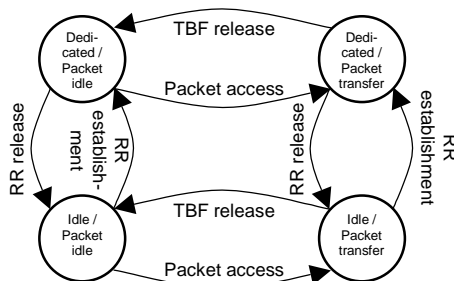


**Figure 3a: RR operating modes and transitions for class A (DTM not supported)**

Figure 3b shows the RR modes and transitions for an MS in modes of operation A (when it supports DTM) and B. In the mode of operation B there are three RR modes:

- (Packet) idle mode

- Packet transfer mode

- Dedicated mode (see 04.18)

For a mobile station that supports DTM class A mode of operation, there is an additional RR mode: dual transfer mode. This mode can only be entered via a packet request procedure while in dedicated mode (see 04.18).
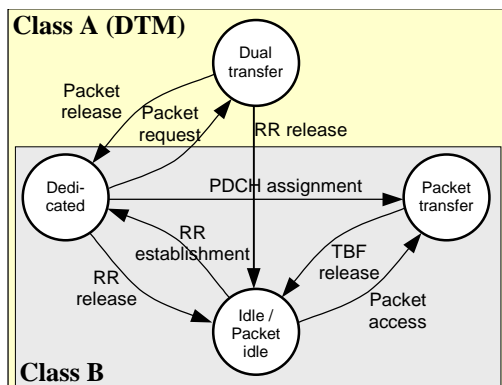


**Figure 3b: RR operating modes and transitions for classes A (DTM supported) and B**

Figure 3c shows the RR modes and transitions for an MS in mode of operation C. The MS can only be attached to either GSM or GPRS:

- when it is GSM attached (and GPRS detached), there are two RR modes: idle mode and dedicated mode

- when it is GPRS attached (and GSM detached), there are two RR modes: packet idle mode and packet transfer mode.
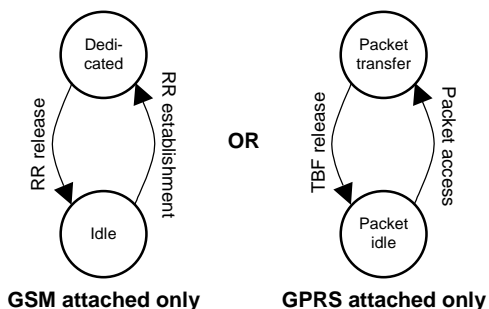


**Figure 3c: RR operating modes and transitions for class C**

# 6.3 Layered overview of radio interface

The GPRS radio interface can be modelled as a hierarchy of logical layers with specific functions. An example of such layering is shown in Figure 3d. The various layers are briefly described in the following subclauses.

The physical layer has been separated into two distinct sub-layers defined by their functions:

- Physical RF layer performs the modulation of the physical waveforms based on the sequence of bits received from the Physical Link layer. The Physical RF layer also demodulates received waveforms into a sequence of bits which are transferred to the Physical Link layer for interpretation.

- Physical Link layer provides services for information transfer over a physical channel between the MS and the Network. These functions include data unit framing, data coding, and the detection and correction of physical medium transmission errors. The Physical Link layer uses the services of the Physical RF layer.

The lower part of the data link layer is defined by following functions:

- The RLC/MAC layer provides services for information transfer over the physical layer of the GPRS radio interface. These functions include backward error correction procedures enabled by the selective retransmission of erroneous blocks. The MAC function arbitrates access to the shared medium between a multitude of MSs and the Network. The RLC/MAC layer uses the services of the Physical Link layer. The layer above RLC/MAC (i.e.,

LLC described in GSM 03.60 [3] and defined in GSM 04.64 [8]) uses the services of the RLC/MAC layer on the Um interface.
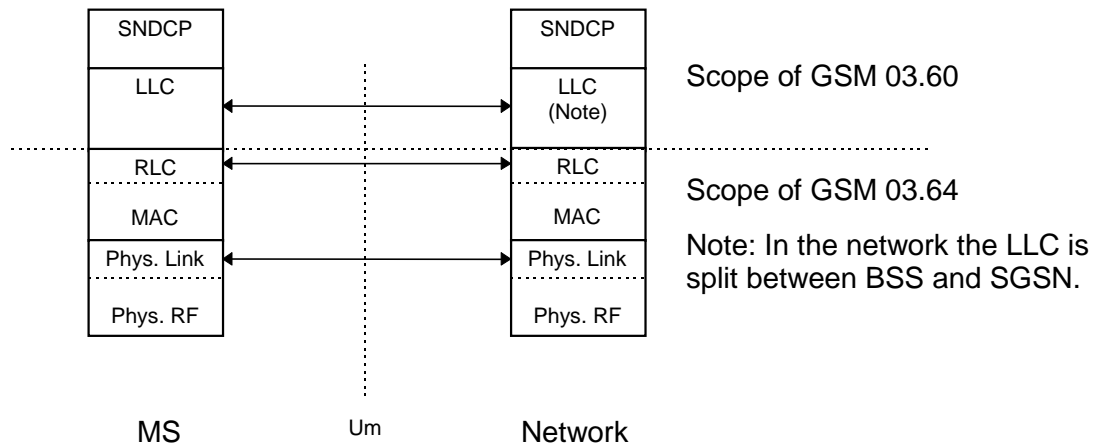


**Figure 3d: GPRS MS – Network Reference Model**

# 6.4 Physical RF Layer

The GSM Physical RF layer is defined in GSM 05.xx series recommendations, which specify among other things:

- The carrier frequencies characteristics and GSM radio channel structures (GSM 05.02 [11]);

- The modulation of the transmitted wave forms and the raw data rates of GSM channels (GSM 05.04 [13]); and

- The transmitter and receiver characteristics and performance requirements (GSM 05.05 [14]).

In the case of EGPRS, the modulation format is inherently signalled by the rotation factor of the training sequences as specified in GSM 05.04 [13] and GSM 05.02 [11] , enabling blind detection in the receiver.

# 6.5 Physical Link Layer

The Physical Link layer operates above the physical RF layer to provide a physical channel between the MS and the Network.

## 6.5.1 Layer Services

The purpose of the Physical Link layer is to convey information across the GSM radio interface, including RLC/MAC information. The Physical Link layer supports multiple MSs sharing a single physical channel.

The Physical Link layer provides communication between MSs and the Network.

The Physical Link layer control functions provide the services necessary to maintain communications capability over the physical radio channel between the Network and MSs. Radio subsystem link control procedures are currently specified in GSM 05.08 [15]. Network controlled handovers are not used in the GPRS service. MS performed cell-reselection is used, see subclause 6.5.6.

## 6.5.2 Layer Functions

The Physical Link layer is responsible for:

- Forward Error Correction (FEC) coding, allowing the detection and correction of transmitted code words and the indication of uncorrectable code words. The coding schemes are described in subclause 6.5.5.

- Interleaving of one Radio Block over four bursts in consecutive TDMA frames, as specified in GSM 05.03 [12].

- Procedures for detecting physical link congestion.

The Physical Link layer control functions include:

- Synchronisation procedures, including means for determining and adjusting the MS Timing Advance to correct for variances in propagation delay , GSM 05.10 [16];

- Monitoring and evaluation procedures for radio link signal quality;

- Cell (re-)selection procedures;

- Transmitter power control procedures; and

- Battery power conservation procedures, e.g. Discontinuous Reception (DRX) procedures.

## 6.5.3 Service Primitives

Table 2 lists the service primitives provided by the Physical Link layer to RLC/MAC layer. More detailed description is given in GSM 04.04 [4].

**Table 2: Service primitives provided by the Physical link layer**

| Name | Request | indication | response | confirm | Comments |
|------|---------|-----------|----------|---------|----------|
| PH-DATA | X | X | | | Used to pass message units containing frames used for RLC/MAC layer respective peer-to-peer communications to and from the physical layer. |
| PH-RANDOM ACCESS | X | X | | X | Used to request and confirm (in the MS) the sending of a random access frame and to indicate (in the network) the arrival of a random access frame. |
| PH-CONNECT | | X | | | Used to indicate that the physical connection on the packet data physical channel has been established. |
| PH-READY-TO-SEND | X | | | | Used by the physical layer to trigger, if applicable, piggy backing, the start of timer for the RLC/MAC layer and the forwarding a data unit to the physical layer |
| PH-EMPTY-FRAME | X | | | | Used by the RLC/MAC layer to indicate that no frame has to be transmitted after receiving the PH-READY-TO-SEND primitive |

## 6.5.4 Radio Block Structure

Different Radio Block structures for data transfer and control message transfer purposes are defined. The Radio Block structure for data transfer is different for GPRS and EGPRS, whereas the same Radio Block structure is used for control messages. For detailed definition of radio block structure, see GSM 04.60 [7].

For GPRS, a Radio Block for data transfer consists of one MAC Header, one RLC header and one RLC Data Block. It is always carried by four normal bursts.

| Radio Block | | | |
|------|------|------|------|
| MAC header | RLC header | RLC data | BCS |

**Figure 4: Radio Block structure for data transfer for GPRS**

The MAC header contains control fields which are different for uplink and downlink directions. The MAC header has constant length, 8 bits.

The RLC header contains control fields which are different for uplink and downlink directions. The RLC header has variable length.

The RLC data field contains octets from one or more LLC PDUs.

The Block Check Sequence (BCS) is used for error detection.

For EGPRS, a Radio Block for data transfer consists of one RLC/MAC header and one or two RLC Data Blocks. It is always carried by four normal bursts. The interleaving depends on the MCS used.

| Radio Block | | | | |
|---|---|---|---|---|
| RLC/MAC header | HCS | RLC data | | BCS |

**Figure 5: Radio Block structure for data transfer for EGPRS**

The RLC/MAC header contains control fields which are different for uplink and downlink directions. The RLC/MAC header has variable length.

The RLC data field contains octets from one or more LLC PDUs.

The Block Check Sequence (BCS) is used for error detection of the data part.

The Header Check Sequence (HCS) is used for error detection of the header part.

The header part is independently coded from the data part and has its own check sequence. Tail biting (i.e. no explicit tail bits are appended before encoding. The encoder is initialised with the last information bits enabling tail biting decoding in the receiver) is used to reduce the size of the header.

For GPRS and EGPRS, a Radio Block for control message transfer consists of one MAC header and one RLC/MAC Control Block. It is always carried by four normal bursts.

| Radio Block | | |
|---|---|---|
| MAC header | RLC/MAC Control Message | BCS |

**Figure 6: Radio Block structure for control message for GPRS and EGPRS**

The MAC header contains control fields which are different for uplink and downlink directions. The MAC header has constant length, 8 bits.

The Block Check Sequence (BCS) is used for error detection.

The RLC/MAC Control message field contains one RLC/MAC control message.

# 6.5.5 Channel Coding

NOTE: The text in this subclause is informative. The normative text is in GSM 05.03 [12]. Where there is a conflict between these descriptions, the normative text has precedence.

Four coding schemes, CS-1 to CS-4, are defined for the GPRS packet data traffic channels. For all other GPRS packet control channels than Packet Random Access Channel (PRACH) and Packet Timing Advance Control Channel on Uplink (PTCCH/U), coding scheme CS-1 is always used. For access bursts on PRACH, two coding schemes are specified.

All coding schemes (CS-1 to CS-4) are mandatory for MSs supporting GPRS. CS-1 is mandatory for a network supporting GPRS.

Nine modulation and coding schemes, MCS-1 to MCS-9, are defined for the EGPRS packet data traffic channels. For all EGPRS packet control channels the corresponding GPRS control channel coding is used. MSs supporting EGPRS shall support MCS-1 to MCS-9 in downlink and MCS-1 to MCS-4 in uplink. In case an MS supporting EGPRS is 8-PSK capable in uplink, it shall also support MCS-5 to MCS-9 in uplink. A network supporting EGPRS may support only some of the MCSs.

## 6.5.5.1 Channel coding for PDTCH

### 6.5.5.1.1 Channel coding for GPRS PDTCH

Four different coding schemes, CS-1 to CS-4, are defined for the GPRS Radio Blocks carrying RLC data blocks. The block structures of the coding schemes are shown in Figure 7 and Figure 8.
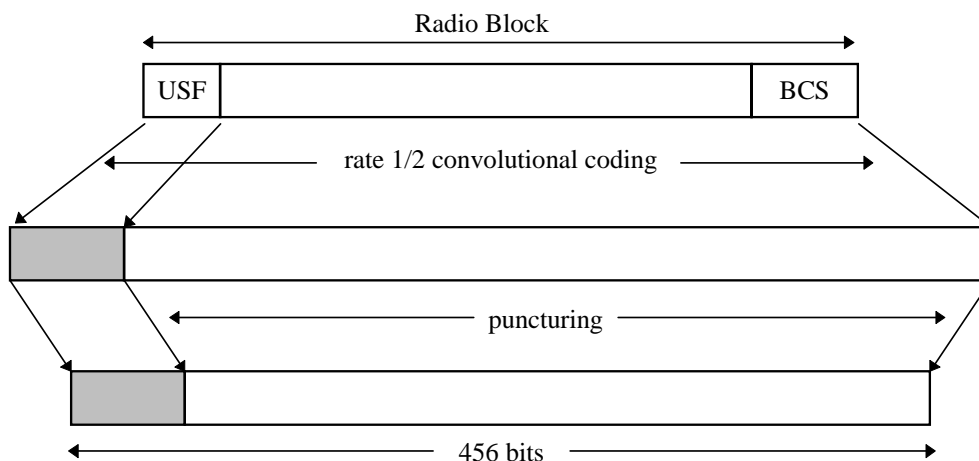
**Figure 7: Radio Block structure for CS-1 to CS-3**
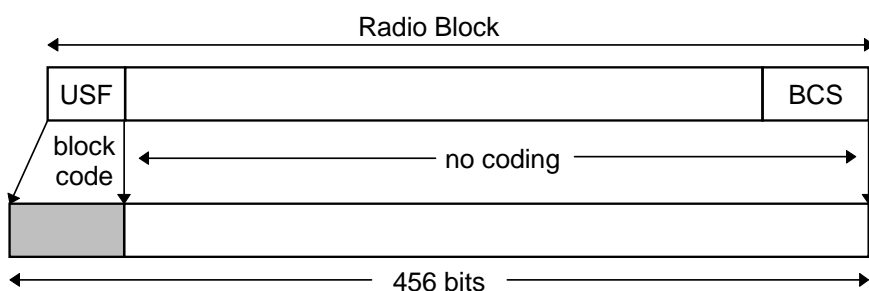


**Figure 8: Radio Block structure for CS-4**

The first step of the coding procedure is to add a Block Check Sequence (BCS) for error detection.

For CS-1 - CS-3, the second step consists of pre-coding USF (except for CS-1), adding four tail bits and a half rate convolutional coding for error correction that is punctured to give the desired coding rate.

For CS-4 there is no coding for error correction.

The details of the codes are shown in table 3, including:

- the length of each field;

- the number of coded bits (after adding tail bits and convolutional coding);

- the number of punctured bits;

- the data rate, including the RLC header and RLC information.

**Table 3: Coding parameters for the GPRS coding schemes.**

| Scheme | Code rate | USF | Pre-coded USF | Radio Block excl. USF and BCS | BCS | Tail | Coded bits | Punctured bits | Data rate kb/s |
|--------|-----------|-----|---------------|-------------------------------|-----|------|------------|----------------|----------------|
| CS-1 | 1/2 | 3 | 3 | 181 | 40 | 4 | 456 | 0 | 9.05 |
| CS-2 | ≈2/3 | 3 | 6 | 268 | 16 | 4 | 588 | 132 | 13.4 |
| CS-3 | ≈3/4 | 3 | 6 | 312 | 16 | 4 | 676 | 220 | 15.6 |
| CS-4 | 1 | 3 | 12 | 428 | 16 | - | 456 | - | 21.4 |

CS-1 is the same coding scheme as specified for SACCH in GSM 05.03 [12]. It consists of a half rate convolutional code for FEC and a 40 bit FIRE code for BCS (and optionally FEC).

CS-2 and CS-3 are punctured versions of the same half rate convolutional code as CS-1 for FEC.

CS-4 has no FEC.

CS-2 to CS-4 use the same 16 bit CRC for BCS. The CRC is calculated over the whole uncoded RLC Data Block including MAC Header.

The USF has 8 states, which are represented by a binary 3 bit field in the MAC Header.

For CS-1, the whole Radio Block is convolutionally coded and USF needs to be decoded as part of the data.

All other coding schemes generate the same 12 bit code for USF. The USF can be decoded either as a block code or as part of the data.

In order to simplify the decoding, the stealing bits (defined in GSM 05.03 [12] ) of the block are used to indicate the actual coding scheme.

## 6.5.5.1.2 Channel coding for EGPRS PDTCH

Nine different modulation and coding schemes, MCS-1 to MCS-9, are defined for the EGPRS Radio Blocks (4 bursts, 20ms) carrying RLC data blocks. The block structures of the coding schemes are shown from Figure 10 to Figure 18 and in Table 4. A general description of the MCSs is given in Figure 9.

The MCSs are divided into different families A, B and C. Each family has a different basic unit of payload: 37 (and 34), 28 and 22 octets respectively. Different code rates within a family are achieved by transmitting a different number of payload units within one Radio Block. For families A and B, 1, 2 or 4 payload units are transmitted, for family C, only 1 or 2 payload units are transmitted.

When 4 payload units are transmitted (MCS-7, MCS-8 and MCS-9), these are splitted into two separate RLC blocks (i.e. with separate sequence numbers and BCSs). These blocks in turn are interleaved over two bursts only, for MCS-8 and MCS-9. For MCS-7, these blocks are interleaved over four bursts. All the other MCSs carry one RLC block which is interleaved over four bursts. When switching to MCS-3 or MCS-6 from MCS-8, 6 padding octets are added to the data octets.
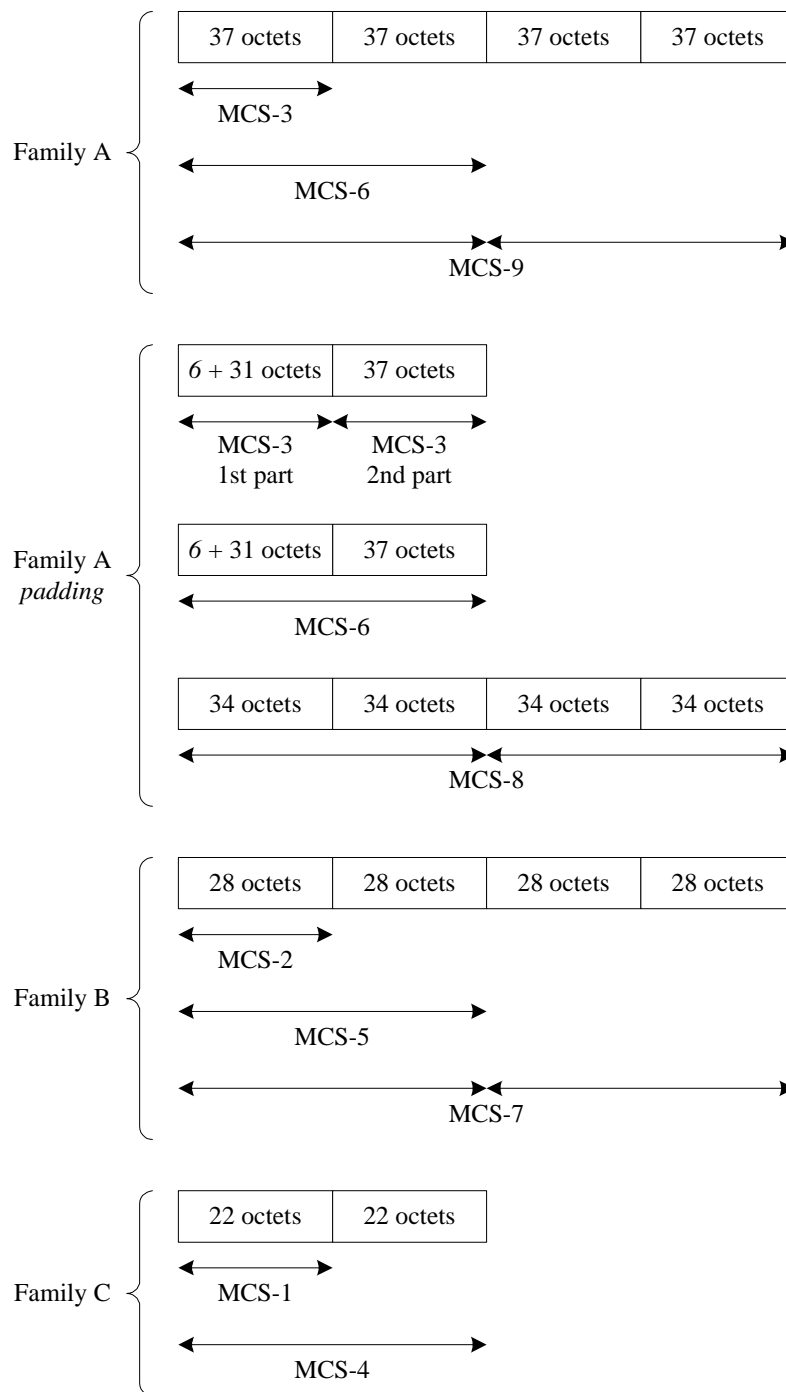
**Figure 9: General description of the Modulation and Coding Schemes for EGPRS**

To ensure strong header protection, the header part of the Radio Block is independently coded from the data part of the Radio Block (8 bit CRC calculated over the header -excl. USF- for error detection, followed by rate 1/3 convolutional coding –and eventually puncturing- for error correction). Three different header formats are used, one for MCS-7, MCS-8 and MCS-9, one for MCS-5 and MCS-6 and one for MCS-1 to MCS-4. The two first formats are for 8PSK modes, the difference being in the number of Sequence Numbers carried (2 for MCS-7, -8 and -9, 1 for MCS-5 and –6). The third format is common to all GMSK modes. The header is always interleaved over four bursts. See 3GPP TS 04.60 [7] for more details.

Following figures show the coding and puncturing for all the Modulation and Coding Schemes, for downlink traffic.
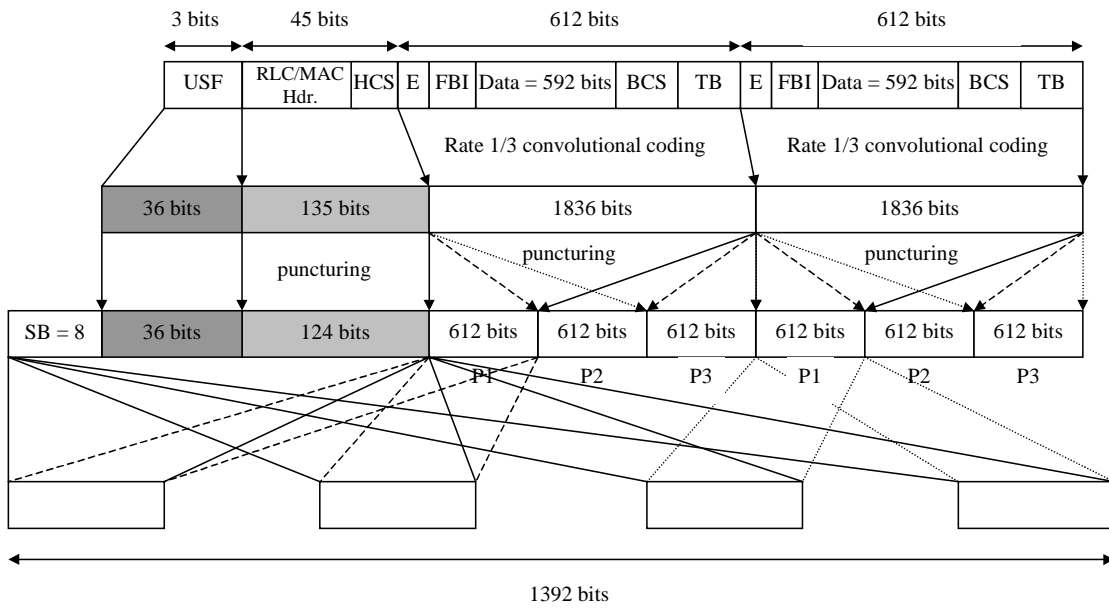
**Figure 10: Coding and puncturing for MCS-9; uncoded 8PSK, two RLC blocks per 20ms**
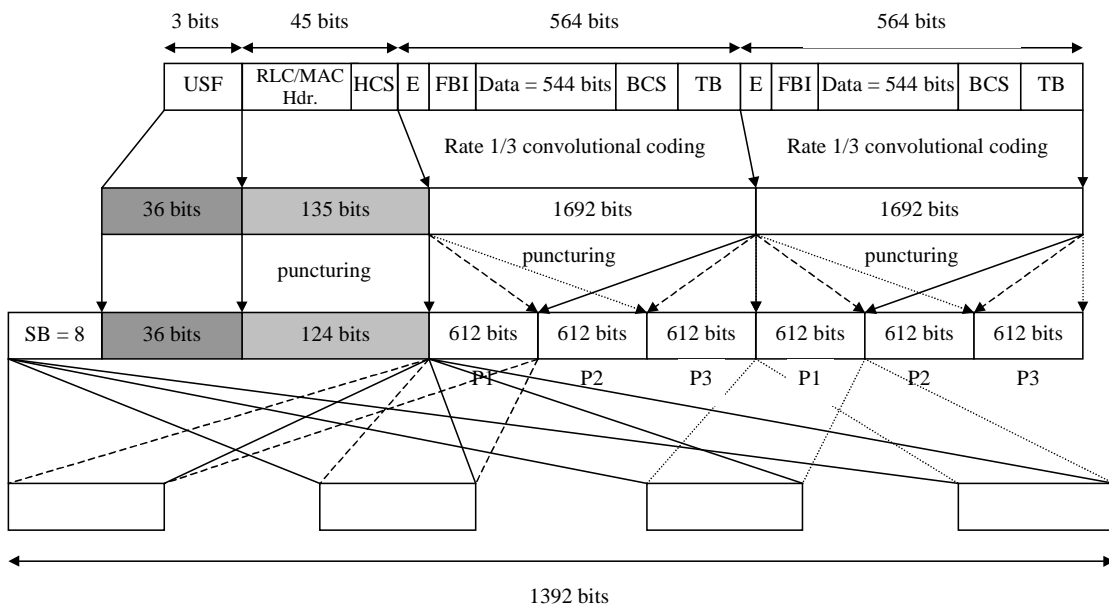


**Figure 11: Coding and puncturing for MCS-8; rate 0.92 8PSK, two RLC blocks per 20ms**
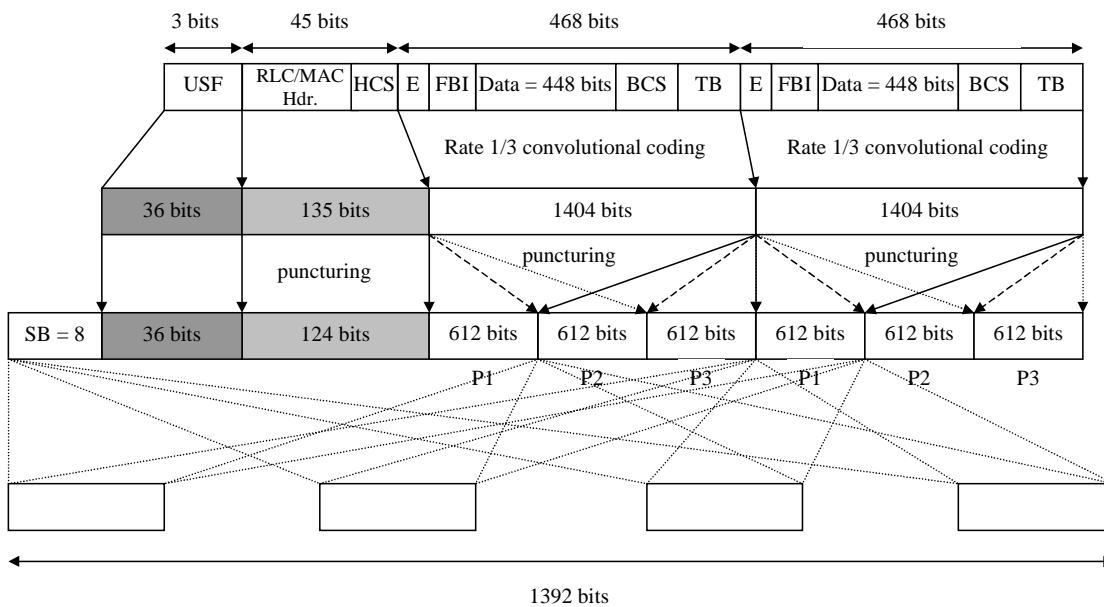
**Figure 12: Coding and puncturing for MCS-7; rate 0.76 8PSK, two RLC blocks per 20ms**
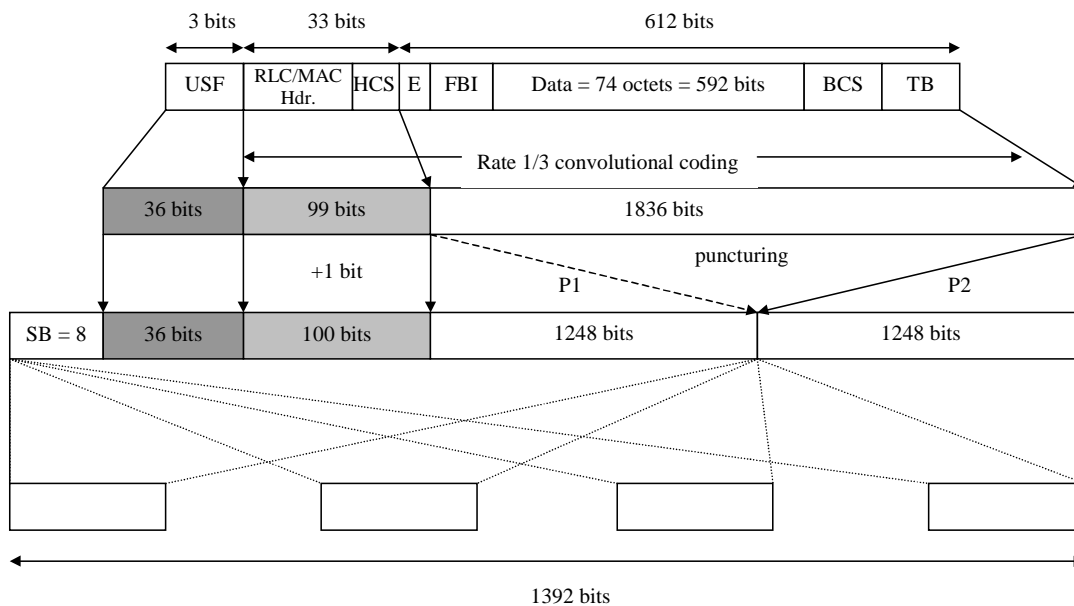


**Figure 13: Coding and puncturing for MCS-6; rate 0.49 8PSK, one RLC block per 20 ms**
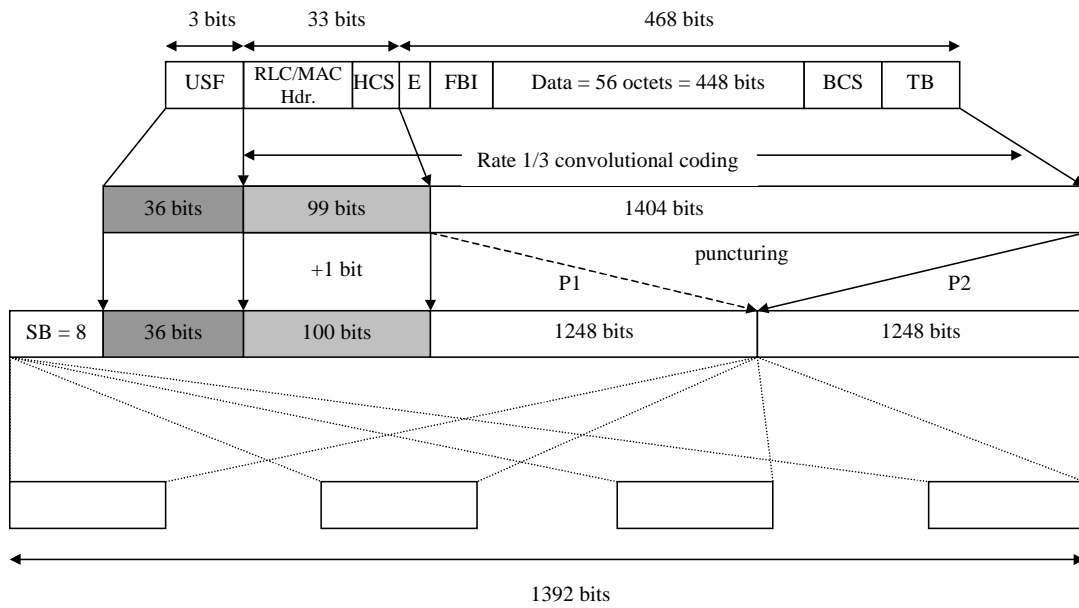
**Figure 14: Coding and puncturing for MCS-5; rate 0.37 8PSK, one RLC block per 20 ms**
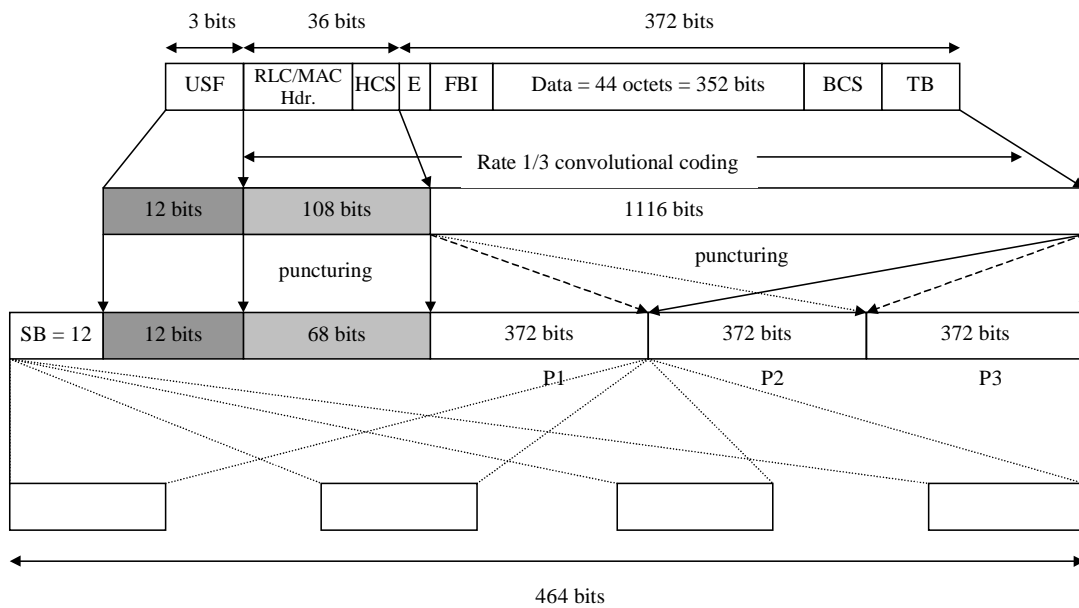


**Figure 15: Coding and puncturing for MCS-4; uncoded GMSK, one RLC block per 20 ms**
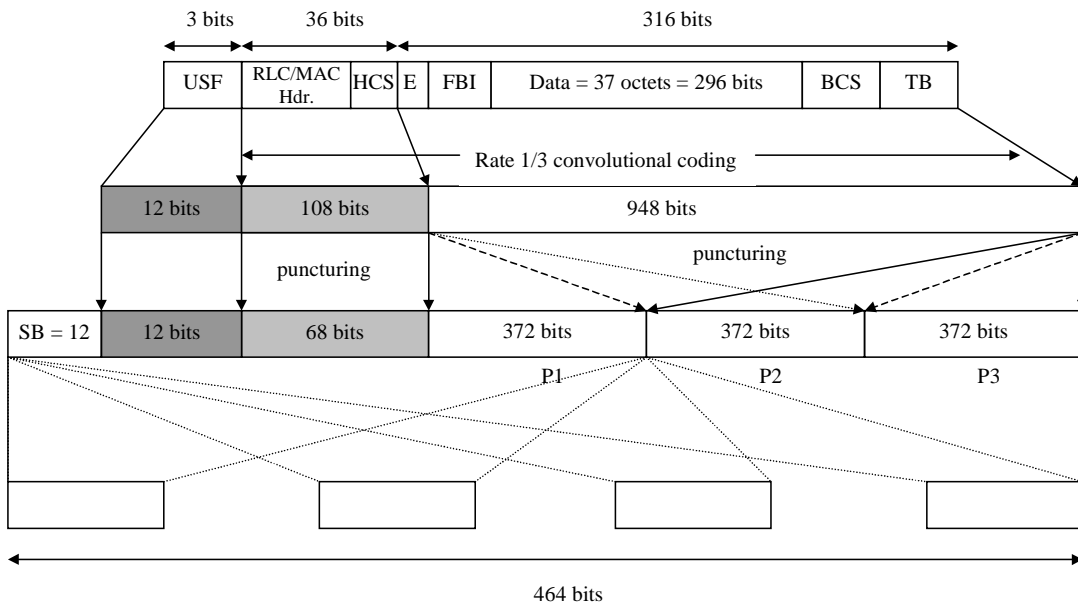
**Figure 16: Coding and puncturing for MCS-3; rate 0.85 GMSK, one RLC block per 20 ms**
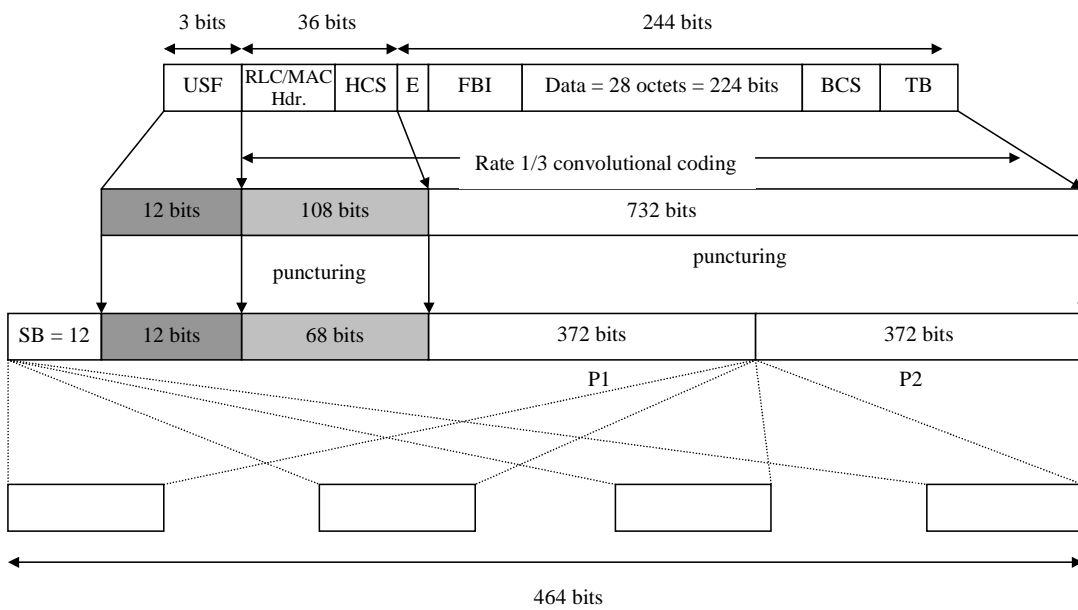


**Figure 17: Coding and puncturing for MCS-2; rate 0.66 GMSK, one RLC block per 20 ms**
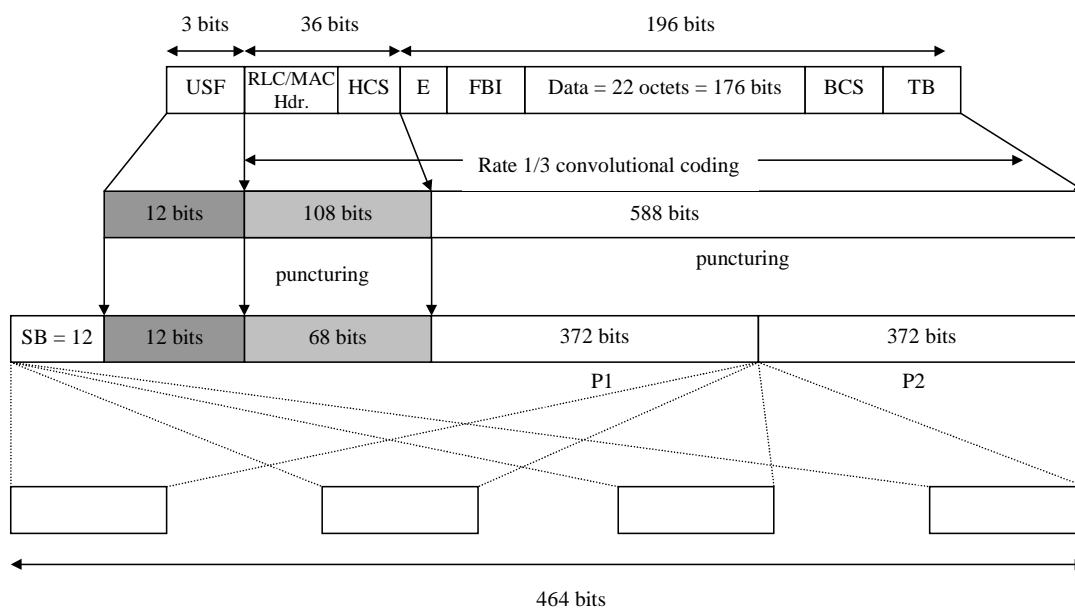
**Figure 18: Coding and puncturing for MCS-1; rate 0.53 GMSK, one RLC block per 20 ms**

The USF has 8 states, which are represented by a binary 3 bit field in the MAC Header. The USF is encoded to 12 symbols similarly to GPRS, (i.e., 12 bits for GMSK modes and 36 bits for 8PSK modes).

The FBI (Final Block Indicator) bit and the E (Extension) bit defined in 3GPP TS 04.60 [7] do not require extra protection: they are encoded along with the data part.

The first step of the coding procedure is to add a Block Check Sequence (BCS) for error detection.

The second step consists of adding six tail bits (TB) and a 1/3 rate convolutional coding for error correction that is punctured to give the desired coding rate. The $P_i$ for each MCS correspond to different puncturing schemes achieving the same coding rate. The bits indicating the MCS used are in the coded header. In both 8PSK and GMSK modes the stealing bits (SB) of the block are used to indicate the header formats. There are eight SB for 8PSK mode which allow to indicate four header formats. There are twelve SB for GMSK mode which allow to indicate two header formats: the first eight of the twelve SB indicate CS-4.

The details of the EGPRS coding schemes are shown in table 4:

**Table 4: Coding parameters for the EGPRS coding schemes**

| Scheme | Code rate | Header Code rate | Modulation | RLC blocks per Radio Block (20ms) | Raw Data within one Radio Block | Family | BCS | Tail payload | HCS | Data rate kb/s |
|--------|-----------|------------------|------------|-----------------------------------|---------------------------------|--------|-----|--------------|-----|----------------|
| MCS-9 | 1.0 | 0.36 | 8PSK | 2 | 2x592 | A | 2x12 | 2x6 | 8 | 59.2 |
| MCS-8 | 0.92 | 0.36 | | 2 | 2x544 | A | | | | 54.4 |
| MCS-7 | 0.76 | 0.36 | | 2 | 2x448 | B | | | | 44.8 |
| MCS-6 | 0.49 | 1/3 | | 1 | 592 *48*+544 | A | 12 | 6 | | 29.6 27.2 |
| MCS-5 | 0.37 | 1/3 | | 1 | 448 | B | | | | 22.4 |
| MCS-4 | 1.0 | 0.53 | GMSK | 1 | 352 | C | | | | 17.6 |
| MCS-3 | 0.85 | 0.53 | | 1 | 296 *48*+248 and 296 | A | | | | 14.8 13.6 |
| MCS-2 | 0.66 | 0.53 | | 1 | 224 | B | | | | 11.2 |
| MCS-1 | 0.53 | 0.53 | | 1 | 176 | C | | | | 8.8 |

NOTE:  The italic captions indicate the 6 octets of padding when retransmitting an MCS-8 block with MCS-3 or MCS-6. For MCS-3, the 6 octets of padding are sent every second block (see 3GPP TS 04.60).

## 6.5.5.2 Channel coding for PACCH, PBCCH, PAGCH, PPCH,PNCH and PTCCH

The channel coding for the PACCH, PBCCH, PAGCH, PPCH,PNCH and downlink PTCCH is the same as the coding scheme CS-1 presented in subclause 6.5.5.1.

The coding scheme used for uplink PTCCH is the same as for PRACH.

## 6.5.5.2a Channel coding for CPBCCH, CPAGCH, CPPCH, CPNCH, and CSCH

The channel coding for the CPBCCH, CPAGCH, CPPCH, and CPNCH is the same as the coding scheme CS-1 presented in subclause 6.5.5.1. The channel coding for the CSCH is identical to SCH.

## 6.5.5.3 Channel Coding for the PRACH and CPRACH

Two types of packet access burst may be transmitted on the PRACH: an 8 information bits access burst or an 11 information bits access burst called the extended packet access burst. The mobile shall support both access bursts. The channel coding for both burst formats is indicated in the following subclauses. Only the 11 information bits access burst may be transmitted on the CPRACH.

### 6.5.5.3.1 Coding of the 8 data bit Packet Access Burst

The channel coding used for the burst carrying the 8 data bit packet access uplink message is identical to the coding of the access burst as defined for random access channel in GSM 05.03 [12].

### 6.5.5.3.2 Coding of the 11 data bit Packet Access Burst

The channel coding for 11 bit access burst is the punctured version of the same coding as used for 8 bit access burst.

## 6.5.6 Cell Re-selection

NOTE:  The text in this subclause is informative. The normative text is in GSM 03.22 and GSM 05.08 [15]. Where there is a conflict between these descriptions, the normative text has precedence.

In GPRS Packet Idle and Packet Transfer modes, cell re-selection is performed by the MS, except for a class A MS (see GSM 02.60 [2] ) while in dedicated mode in which case the cell is determined by the network according to the handover procedures.

The new cell re-selection criteria C31 and C32 are provided as a complement to the current GSM cell re-selection criteria. This provides a more general tool to make cell planning for GPRS as similar to existing planning in GSM as possible.C31 is a signal strength criterion used to decide whether prioritised cell re-selection shall be used. For cells that fulfil the C31 criterion, the cell with highest priority class shall be selected. If more than one cell has the highest priority, the one of those with the highest C32 value shall be selected. If no cell fulfils the C31 criterion, the one among all cells with the highest C32 value shall be selected.

C32 is an improvement of C2. It applies an individual offset and hysteresis value to each pair of cells, as well as the same temporary offsets as for C2. Additional hysteresis values apply for a cell re-selection that requires cell or routing area update.

Cell re-selection procedure apply to the MSs attached to GPRS if a PBCCH exists in the serving cell. If the PBCCH is not allocated, then the MS shall perform cell re-selection according to the C2 criteria.

In addition, the network may control the cell re-selection as described in subclause 6.5.6.3.

## 6.5.6.1 Measurements for Cell Re-selection

The MS shall measure the received RF signal strength on the BCCH frequencies of the serving cell and the neighbour cells as indicated in the BA-GPRS list, and calculate the received level average (RLA) for each frequency, as specified in GSM 05.08 [15]. In addition the MS shall verify the BSIC of the cells. Only channels with the same BSIC as broadcast together with BA-GPRS on PBCCH shall be considered for re-selection.
A COMPACT capable MS shall in addition perform the above tasks for any CPBCCH, either transmitted in the serving cell or indicated, by way of frequency and time group in the BA-GPRS list.
Any cell having a CPBCCH indicated in the BA-GPRS list shall be time synchronized to that cell, as specified in GSM 05.10 [16].
When the number of downlink PDCHs assigned to certain types of multislot MS (see GSM 05.02 [11] , annex B) does not allow them to perform measurements within the TDMA frame, the network shall provide measurement windows to ensure that the MS can perform a required number of measurements. The network shall provide periods of inactivity during a fixed allocation to allow the MS to make adjacent cell power measurements and BSIC detection.

## 6.5.6.2 Broadcast Information

The PBCCH broadcasts GPRS specific cell re-selection parameters for serving and neighbour cells, including the BA (GPRS) list. A BA (GPRS) identifies the neighbour cells, including BSIC, that shall be considered for GPRS cell (re-selection (not necessary the same as for GSM in Idle or circuit switched mode)).

## 6.5.6.3 Optional measurement reports and network controlled cell re-selection

It shall be possible for the network to order the mobile stations to send measurement reports to the network and to suspend its normal cell re-selection, and instead to accept decisions from the network. This applies to both Packet idle mode and Packet transfer mode.

The degree to which the mobile station shall resign its radio network control shall be variable, and be ordered in detail by the parameter NETWORK_CONTROL_ORDER.

Two sets of parameters are broadcast on PBCCH and are valid in Packet transfer and Packet idle modes respectively. NETWORK_CONTROL_ORDER can also be sent individually to an MS on PACCH, in which case it overrides the broadcast parameter.

Additionally, the network may request extended measurement reports from the MS and the reporting shall be maintained in packet idle mode. The reports may include interference measurements (see subclause 6.5.8.3.2). Measurement reports shall be sent individually from each MS as RLC transmissions.

When a class A mobile station is simultaneously involved in a circuit switched service and in a GPRS transfer, the network controlled cell re-selection procedures (NC1 or NC2 modes of operation) should not be used. In this case, handover for the circuit switched service has precedence over GPRS network controlled cell re-selection, and the MS shall stop sending measurement reports and ignore cell change orders.

## 6.5.7 Timing Advance

NOTE: The text in this subclause is informative. The normative text is in GSM 04.60 [7] and GSM 05.10 [16]. Where there is a conflict between these descriptions, the normative text has precedence.

The timing advance procedure is used to derive the correct value for timing advance that the MS has to use for the uplink transmission of radio blocks. This procedure is not used in dual transfer mode, in which case the timing advance procedures for dedicated mode are used.

The timing advance procedure comprises two parts:

- initial timing advance estimation;

- continuous timing advance update.

### 6.5.7.1 Initial timing advance estimation

The initial timing advance estimation is based on the single access burst carrying the Packet Channel Request. The Packet Uplink Assignment or Packet Downlink Assignment then carries the estimated timing advance value to the MS. This value shall be used by the MS for the uplink transmissions until the continuous timing advance update provides a new value (see subclause 6.5.7.2.). Two special cases exist:

- when Packet Queuing Notification is used the initial estimated timing advance may become too old to be sent in the Packet Downlink (/Uplink) Assignment

- when Packet Downlink (/Uplink) Assignment is to be sent without prior paging (i.e., in the Ready state), no valid timing advance value may be available.

Then the network has three options:

- Packet Polling Request can then be used to trigger the transmission of Packet Control Acknowledgement. This message can be formatted as four access burst from which the timing advance can be estimated.

- Packet Downlink (/Uplink) Assignment can be sent without timing advance information. In that case it is indicated to the MS that it can only start the uplink transmission after the timing advance is obtained by the continuous timing advance update procedure.

- The poll bit in the Packet Downlink (/Uplink) Assignment message can be set to trigger the transmission of Packet Control Acknowledgement. This can be used if System information indicates that acknowledgement is access bursts.

For the case where timing advance information is not provided in the assignment message, the mobile is not allowed to send normal bursts on the uplink until it receives a valid timing advance either in Packet Timing Advance/Power Control message or through the continuous timing advance procedure.

### 6.5.7.2 Continuous timing advance update

MS in Packet transfer mode shall use the continuous timing advance update procedure. The continuous timing advance update procedure is carried on the PTCCH allocated to the MS.

For uplink packet transfer, within the Packet Uplink Assignment, the MS is assigned Timing Advance Index (TAI) and the PTCCH.

For downlink packet transfer, within the Packet Downlink Assignment, the MS is assigned Timing Advance Index (TAI) and the PTCCH.

The TAI specifies the PTCCH sub-channel used by the MS.

On the uplink, the MS shall send in the assigned PTCCH access burst, which is used by the network to derive the timing advance.

The network analyses the received access burst and determines new timing advance values for all MSs performing the continuous timing advance update procedure on that PDCH. The new timing advance values shall be sent via a

downlink signalling message (TA-message) on PTCCH/D. Network can send timing advance information also in Packet Timing Advance/Power Control and Packet Uplink Ack/Nack messages on PACCH.

## 6.5.7.2.1 Mapping on the multiframe structure

Figure 19 shows the mapping of the uplink access bursts and downlink TA-messages on groups of eight 52-multiframes:

- the TAI value shows the position where a slot is reserved for a MS to send an access burst (e.g. T1 means 52-multiframe number n and idle slot number 2). TAI value defines the used PTCCH sub-channel.

- every second PDCH multiframe starts a downlink TA-message.

52-multiframe number n:

| uplink | | | TAI=0 | | | | | | | TAI=1 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 0 | B3 | B4 | B5 | 1 | B6 | B7 | B8 | 2 | B9 | B10 | B11 | 3 |

downlink            TA_message 1                          TA message 1

52-multiframe number n + 1:

| uplink | | | TAI=2 | | | | | | | TAI=3 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 4 | B3 | B4 | B5 | 5 | B6 | B7 | B8 | 6 | B9 | B10 | B11 | 7 |

downlink            TA message 1                          TA message 1

52-multiframe number n + 2:

| uplink | | | TAI=4 | | | | | | | TAI=5 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 8 | B3 | B4 | B5 | 9 | B6 | B7 | B8 | 10 | B9 | B10 | B11 | 11 |

downlink            TA message 2                          TA message 2

52-multiframe number n + 3:

| uplink | | | TAI=6 | | | | | | | TAI=7 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 12 | B3 | B4 | B5 | 13 | B6 | B7 | B8 | 14 | B9 | B10 | B11 | 15 |

downlink            TA message 2                          TA message 2

52-multiframe number n + 4:

| uplink | | | TAI=8 | | | | | | | TAI=9 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 16 | B3 | B4 | B5 | 17 | B6 | B7 | B8 | 18 | B9 | B10 | B11 | 19 |

downlink            TA message 3                          TA message 3

52-multiframe number n + 5:

| uplink | | | TAI=10 | | | | | | | TAI=11 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 20 | B3 | B4 | B5 | 21 | B6 | B7 | B8 | 22 | B9 | B10 | B11 | 23 |

downlink            TA message 3                          TA message 3

52-multiframe number n + 6:

| uplink | | | TAI=12 | | | | | | | TAI=13 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 24 | B3 | B4 | B5 | 25 | B6 | B7 | B8 | 26 | B9 | B10 | B11 | 27 |

downlink            TA message 4                          TA message 4

52-multiframe number n + 7:

| uplink | | | TAI=14 | | | | | | | TAI=15 | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| B0 | B1 | B2 | 28 | B3 | B4 | B5 | 29 | B6 | B7 | B8 | 30 | B9 | B10 | B11 | 31 |

downlink            TA message 4                          TA message 4

B0 - B11 = Radio blocks
Idle frames are numbered from 1 to 31 [odd numbers]
PTCCH frames are numbered from 0 to 30 [even numbers]

**Figure 19: Mapping of the uplink access bursts and downlink timing advance signalling messages**

The BTS shall update the timing advance values in the next TA-message following the access burst. To illustrate this, an MS that transmits an access burst in frames numbered 0, 2, 4, or 6 receives its updated timing advance value in TA message 2. This MS can also find this updated timing advance value in subsequent TA messages 3, 4, and 1, but only has to read these if TA message 2 was not received correctly.

An MS entering the Transfer state shall ignore the TA-messages until the MS has sent its first access burst. This is to avoid the use of timing advance values, derived from access bursts sent by the MS that previously used the same TAI.

## 6.5.8 Power control procedure

Power control shall be supported in order to improve the spectrum efficiency and to reduce the power consumption in the MS.

For the uplink, the MS shall follow a flexible power control algorithm, which the network can optimise through a set of parameters. It can be used for both open loop and closed loop power control.

For the downlink, the power control is performed in the BTS. Therefore, there is no need to specify the actual algorithms, but information about the downlink performance is needed. Therefore the MSs have to transfer Channel Quality Reports to the BTS.Power control is not applicable to point-to-multipoint multicast services.

For the detailed specification of power control see GSM 05.08 [15].

### 6.5.8.1 MS output power

The MS shall calculate the RF output power value, $P_{CH}$, to be used on each individual uplink PDCH assigned to the MS:

$$P_{CH} = \min((\Gamma_0 - \Gamma_{CH} - \alpha * (C + 48), PMAX)$$

where

| | |
|---|---|
| $\Gamma_{CH}$ | is an MS and channel specific power control parameter. It is sent to the MS in any resource assigning message. Further, the network can, at any time during a packet transfer, send new $\Gamma_{CH}$ values to the MS on the downlink PACCH. |
| $\Gamma_0$ | is a frequency band dependent constant. |
| $\alpha \in [0,1]$ | is a system parameter. Its default value is broadcast on the PBCCH. Further, MS and channel specific values can be sent to the MS together with $\Gamma_{CH}$. |
| C | is the received signal level at the MS. |
| PMAX | is the maximum allowed output power in the cell. |

All power values are expressed in dBm.

$P_{CH}$ is not used to determine the output power when accessing the cell on PRACH or RACH , in which case PMAX shall be used.

### 6.5.8.2 BTS output power

The BTS shall use constant power on those PDCH radio blocks which contain PBCCH or which may contain PPCH. This power may be lower than the output power used on BCCH. The difference shall be broadcast on PBCCH.

On the other PDCH radio blocks, downlink power control may be used. Thus, a procedure may be implemented in the network to control the power of the downlink transmission based on the Channel Quality Reports.

The network shall ensure that the output power is sufficient for the MS for which the RLC block is intended as well as the MS(s) for which the USF is intended, and that for each MS in packet transfer mode, at least one downlink RLC block per multiframe is transmitted with an output power that is sufficient for that MS, on a block monitored by that MS.

### 6.5.8.3 Measurements at MS side

A procedure shall be implemented in the MS to monitor periodically the downlink Rx signal level and quality from its serving cell.

#### 6.5.8.3.1 Deriving the C value

This subclause comprises information about how the MS shall derive the C value in the power control equation.

The MS shall periodically measure the received signal strength.

In packet idle mode, the MS shall measure the signal strength of the PCCCH or, if PCCCH is not existing, the BCCH.

In packet transfer mode, the MS shall measure the signal strength on BCCH. The same measurements as for cell re-selection are used (see 6.5.6.1.). Alternatively, if indicated by a broadcast parameter, the MS shall measure the signal strength on one of the PDCHs where the MS receives PACCH. This method is suitable in the case where BCCH is in another frequency band than the used PDCHs. It requires that constant output power is used on all downlink PDCH blocks.

The MS shall measure the signal strength of each radio block monitored by the MS. The C value is achieved by filtering the signal strength with a running average filter. The filtering shall normally be continuous between the packet modes. The different filter parameters for the packet modes are broadcast on PBCCH or, if PBCCH does not exist, on BCCH.

The variance of the received signal level within each block shall also be calculated. The filtered value SIGN_VAR shall be included in the channel quality report.

An MS transferring a packet in the uplink with fixed assignment is not required to make signal strength measurements and shall thus update $P_{CH}$ only when it receives new $\Gamma_{CH}$ values.

#### 6.5.8.3.2 Derivation of Channel Quality Report

The channel quality is measured as the interference signal level during the idle frames of the multiframe, when the serving cell is not transmitting.

In packet transfer mode, the MS shall measure the interference signal strength of all eight channels (slots) on the same carrier as the assigned PDCHs.

In packet idle mode, the MS shall measure the interference signal strength on certain channels which are indicated on the PBCCH or, if PBCCH does not exist, on BCCH. If no channels are indicated the MS shall not perform these measurements.

Some of the idle frames and PTCCH frames shall be used for this, while the others are required for BSIC identification and the timing advance procedure, see subclause 6.5.9.

The MS may not be capable of measuring all eight channels when allocated some configurations of channels. The MS shall measure as many channels as its allocation allows considering its multislot capability.

The slots that the MS measures on can be either idle or used by SACCH or PTCCH, depending on the channel type (TCH or PDCH).. The MS shall therefore, for each slot, take the minimum signal strength of one idle frame and one PTCCH frame. Thus the SACCH frames are avoided (except for a TCH/H with two MSs) and only the interference is measured.

The interference, $\gamma_{CH}$ , is achieved by filtering the measured interference in a running average filter. The filtering shall be continuous between the packet modes for channels measured in both modes. The different filter parameters for the packet modes are broadcast on PBCCH or, if PBCCH does not exist, on BCCH.

In packet transfer mode the MS shall transfer the 8 $\gamma_{CH}$ values and the RXQUAL, SIGN_VAR and C values (see subclause 6.5.8.3.1) to the network in the Channel Quality Report included in the PACKET DOWNLINK ACK/NACK message.

### 6.5.8.4 Measurements at BSS side

A procedure shall be implemented in the BSS to monitor the uplink Rx signal level and quality on each uplink PDCH, active as well as inactive.

The BSS shall also measure the Rx signal level and the quality of a specific MS packet transfer.

## 6.5.9    Scheduling the MS activities during the PTCCH and idle frames

The MS shall use the PTCCH and idle frames of the PDCH multiframe for the following tasks:

- BSIC identification for cell re-selection (6.5.6.1)

- Continuous timing advance procedures (6.5.7.2)

- Interference measurements for power control (6.5.8.3.2)

It is not necessary to exactly specify the scheduling of these tasks.

The PTCCH frames used for timing advance signalling is stated in 6.5.7.2.1. During the frames when the MS receives TA-messages it can also make interference measurements. During the frames when the MS transmits access bursts it may also be possible to make measurements on some channels.

The MS shall schedule the BSIC identification as efficiently as possible, using the remaining PTCCH frames and the idle frames and also considering the requirements for interference measurements. When the MS is synchronised to a BTS, it knows the timing of the SCH. Therefore, only a few certain frames are required for BSIC identification. In those frames it may also be possible to make measurements on some channels. When the MS shall synchronise to a new BTS, it has to prioritise that task. It may then use half of the PTCCH and idle frames, i.e. the same amount as available for circuit switched connections.

The remaining PTCCH and idle frames shall be used for interference measurements.

## 6.5.10    Discontinuous Reception (DRX)

> NOTE:    The text in this subclause is informative. The normative text is in GSM 05.02 [11]. Where there is a conflict between these descriptions, the normative text has precedence.

DRX (sleep mode) shall be supported when the MS is in Packet Idle mode. DRX is independent from MM states Ready and Standby.

Negotiation of DRX parameters is per MS. An MS may choose to use DRX or not together with some operating parameters. The following parameters are established:

- **DRX/non-DRX indicator**
  It indicates whether the MS uses DRX or not.

- **DRX period**
  A conditional parameter for MSs using DRX to determine the right paging group. The DRX period is defined by the parameter SPLIT_PG_CYCLE.

- **Non-DRX timer**
  A conditional parameter for MSs using DRX to determine the time period within which the non-DRX mode is kept after leaving the Transfer state. The support for this feature is optional on the network side and the information about the maximum supported value for the timer in the cell is broadcast on PBCCH.

  An MS in DRX mode is only required to monitor the radio blocks defined by its paging group as defined in GSM 05.02 [11].

Paging group definition based on SPLIT_PG_CYCLE is optional on CCCH for both BTS and MS. If not supported, the definition based on BS_PA_MFRMS  shall be used. The parameters used to define the paging group for GPRS are shown in the Table 5, together with the corresponding GSM parameters. BS_PCC_CHANS is the number of PDCHs containing PCCCH. For Compact, BS_PCC_CHANS is the number of radio frequency channels per cell carrying CPCCCHs including the radio frequency channel carrying the CPBCCH.

An MS in non-DRX mode is required to monitor all the radio blocks where PCCCH or (for Compact) CPCCCH may be mapped on the PDCH defined by its paging group.

When page for circuit-switched services is conveyed on PPCH, it follows the same scheduling principles as the page for packet data. The same is valid for scheduling of resource assignments for downlink packet transfers for MSs in Ready State (i.e. where no paging is performed).

The MS may need to monitor also PNCH or CPNCH in the case of PTM-M services.

NOTE: Paging reorganisation may be supported in the same way as for circuit switched GSM.

**Table 5: Parameters for DRX operation**

| Parameter | GPRS | | Corresponding GSM parameters |
|---|---|---|---|
| | PCCCH | CCCH | CCCH |
| DRX period | SPLIT_PG_CYCLE | BS_PA_MFRMS    *) SPLIT_PG_CYCLE **) | BS_PA_MFRMS |
| Blocks not available for PPCH or CPPCH per multiframe | BS_PAG_BLKS_RES  + BS_PBCCH_BLKS | BS_AG_BLKS_RES | BS_AG_BLKS_RES |
| Number of physical channels containing paging; or for Compact, number of radio frequency channels per cell carrying CPCCCHs including the radio frequency channel carrying the CPBCCH. | BS_PCC_CHANS | BS_CC_CHANS | BS_CC_CHANS |

     \*)     Only when DRX period split is not supported.
    \*\*)     Only when DRX period split is supported.

# 6.6 Medium Access Control and Radio Link Control Layer

The Medium Access Control (MAC) and Radio Link Control (RLC) layer operates above the Physical Link layer in the reference architecture. MAC/RLC layer messages and signalling procedures are defined in GSM 04.60 [7] and GSM 04.08 [6].

## 6.6.1 Layer Services

The MAC function defines the procedures that enable multiple MSs to share a common transmission medium, which may consist of several physical channels. The MAC function provides arbitration between multiple MSs attempting to transmit simultaneously and provides collision avoidance, detection and recovery procedures. The operations of the MAC function may allow a single MS to use several physical channels in parallel.

The RLC function defines the procedures for a bitmap selective retransmission of unsuccessfully delivered RLC Data Blocks.

The RCL/MAC function provides two modes of operation:

- unacknowledged operation; and

- acknowledged operation

## 6.6.2 Layer Functions

The GPRS MAC function is responsible for:

- Providing efficient multiplexing of data and control signalling on both uplink and downlink, the control of which resides on the Network side. On the downlink, multiplexing is controlled by a scheduling mechanism. On the uplink, multiplexing is controlled by medium allocation to individual users (e.g., in response to service request).

- For mobile originated channel access, contention resolution between channel access attempts, including collision detection and recovery.

- For mobile terminated channel access, scheduling of access attempts, including queuing of packet accesses.

- Priority handling.

The GPRS RLC function is responsible for:

- Interface primitives allowing the transfer of Logical Link Control layer PDUs (LLC-PDU) between the LLC layer and the MAC function.

- Segmentation and re-assembly of LLC-PDUs into RLC Data Blocks.

- Backward Error Correction (BEC) procedures enabling the selective retransmission of uncorrectable code words.

- Transmission of code words according to the channel conditions,i.e link adaptation.

    NOTE:    The Block Check Sequence for error detection is provided by the Physical Link Layer.

In EGPRS incremental redundancy (IR) mode, RLC function is also responsible for:

- Storing soft values of the erroneous RLC Data Blocks and combining them with the retransmitted RLC Data blocks.

## 6.6.3    Service Primitives

Table 6 lists the service primitives provided by the RLC/MAC layer to the upper layers:

**Table 6: Service primitives provided by the RLC/MAC layer to the upper layers**

| Name | request | indication | response | confirm | comments |
|------|---------|------------|----------|---------|----------|
| RLC/MAC-DATA | x | x | | | used for the transfer of upper layer PDUs. Acknowledged mode of operation in RLC is used. The upper layer shall be able to request high transmission quality via a primitive parameter. |
| RLC/MAC-UNITDATA | x | x | | | used for the transfer of upper layer PDUs. Unacknowledged mode of operation in RLC is used. |
| RLC/MAC-STATUS | | x | | | used to indicate that an error has occurred on the radio interface. The cause for the failure is indicated. |

## 6.6.4    Model of Operation

Each PDCH is a shared medium between multiple MSs and the Network except in dual transfer mode, in which a PDCH may be dedicated to one MS. Direct communication is possible only between an MS and the network.

The GPRS radio interface consists of asymmetric and independent uplink and downlink channels. The downlink carries transmissions from the network to multiple MSs and does not require contention arbitration. The uplink is shared among multiple MSs and requires contention control procedures.

The allocation of radio resources by the PLMN and the use of these resources by the MSs can be broken down into two parts:

- The PLMN allocates radio resources for the GPRS (uplink and downlink) in a symmetric manner.

- The allocated uplink and downlink radio resources for point-to-point, point-to-multipoint multicast or group call service types are used independently of each other. Dependent allocation of uplink and downlink shall be possible, in order to allow simple MSs to transfer data simultaneously in both directions. Allocation of several PDTCHs for one MS is possible.

The access to the GPRS uplink uses a Slotted-Aloha based reservation protocol.

The Network Protocol Data Units (N-PDU) are segmented into the Subnetwork Protocol Data Units (SN-PDU) by the Subnetwork Dependent Convergence (SNDC) protocol and SN-PDUs are encapsulated into one or several LLC frames.

See GSM 03.60 [3] for information on SNDC and LLC. The details on SNDC can be found in GSM 04.65 [9] and the details on LLC can be found in GSM 04.64 [8]. LLC frames are segmented into RLC Data Blocks. At the RLC/MAC layer, a selective ARQ protocol (including block numbering) between the MS and the Network provides retransmission of erroneous RLC Data Blocks. When a complete LLC frame is successfully transferred across the RLC layer, it is forwarded to the LLC layer.



FH  = Frame Header
FCS = Frame Check Sequence
BH  = Block Header
BCS = Block Check Sequence
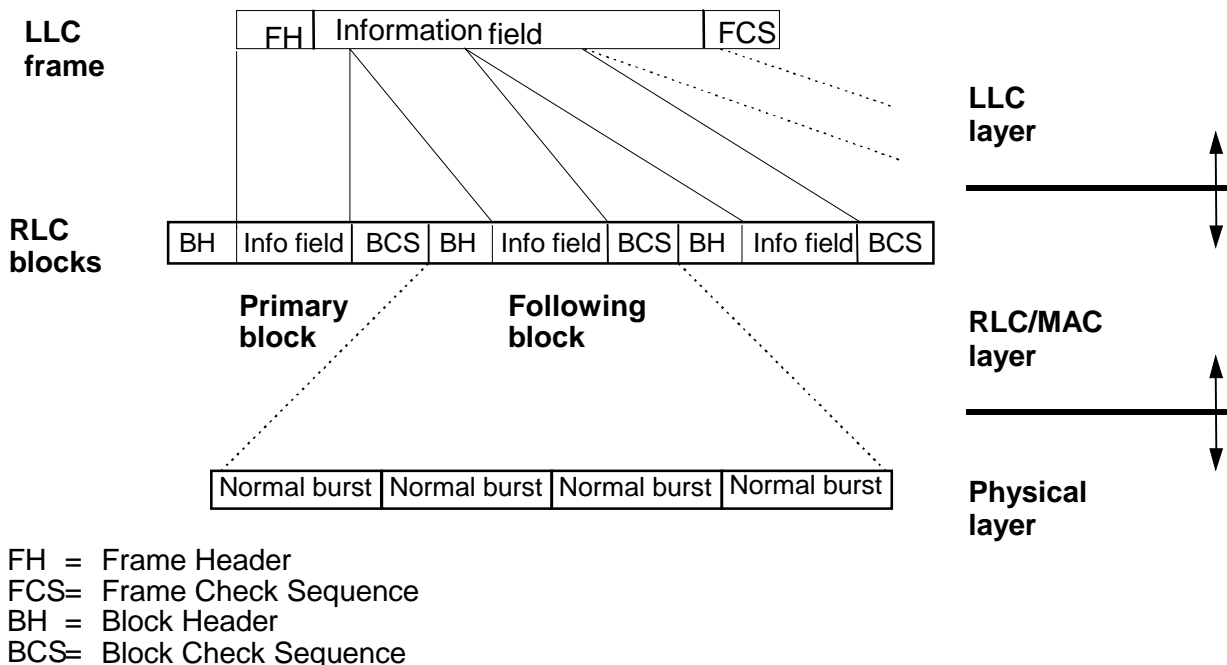
**Figure 20: Transmission and reception data flow for GPRS and EGPRS**

Transmission and reception data flows are same for GPRS and EGPRS, except for  EGPRS MCS-9, MCS-8 and MCS-7, where four normal bursts carry two RLC blocks (one RLC block within two bursts for MCS-9 and MCS-8), as shown on Figures 21 and 22.
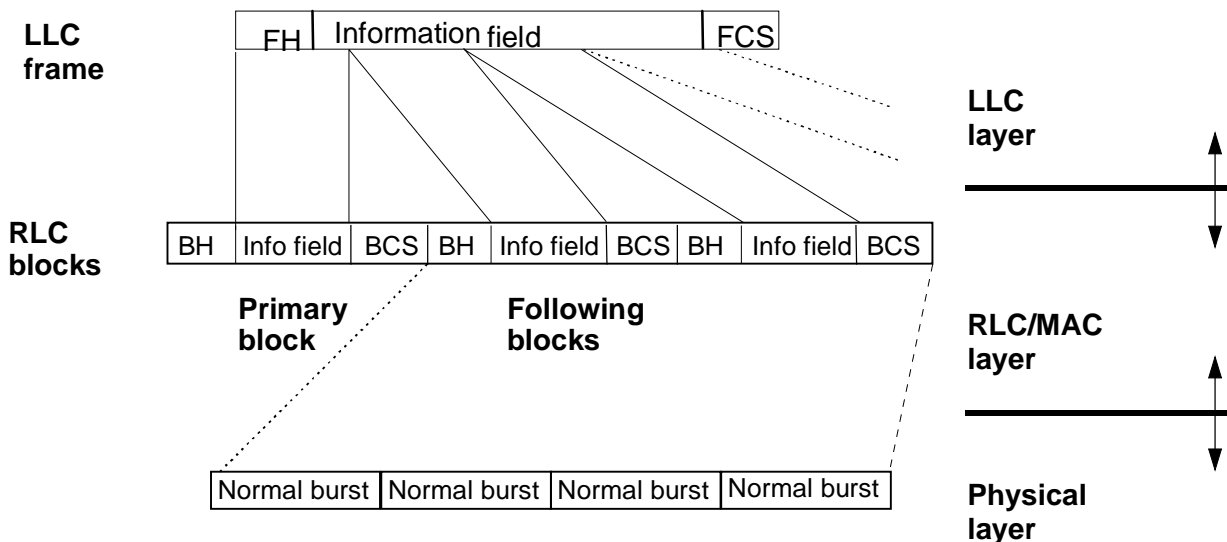


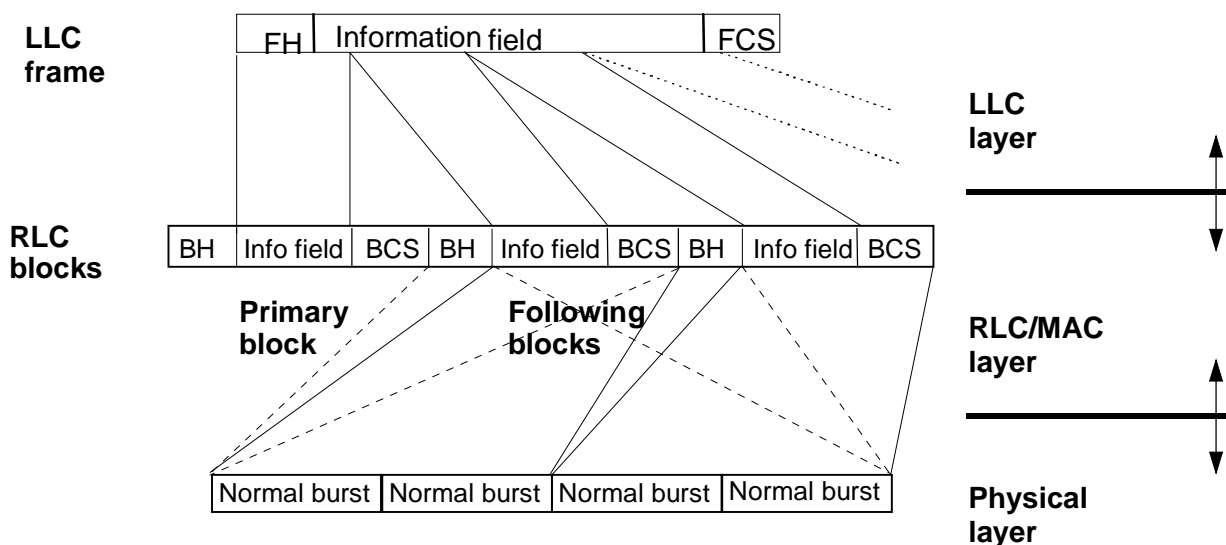**Figure 21: Transmission and reception data flow for EGPRS MCS-7**

**Figure 22: Transmission and reception data flow for EGPRS MCS-9 and MCS-8**

## 6.6.4.1         Multiplexing MSs on the same PDCH

### 6.6.4.1.1         Uplink State Flag: Dynamic Allocation

#### 6.6.4.1.1.1         Multiplexing of GPRS or EGPRS MSs

The Uplink State Flag (USF) is used on PDCH to allow multiplexing of Radio blocks from a number of MSs. USF is used in dynamic and extended dynamic medium access modes. USF is used only in downlink direction.

The USF comprises 3 bits at the beginning of each Radio Block that is sent on the downlink. It enables the coding of 8 different USF states which are used to multiplex the uplink traffic.

On PCCCH, one USF value is used to denote PRACH. The other USF values are used to reserve the uplink for different MSs. On PDCHs not carrying PCCCH, the eight USF values are used to reserve the uplink for different MSs. One USF value shall be used to prevent collision on uplink channel, when MS without USF is using uplink channel. The USF points either to the next uplink Radio Block or the sequence of 4 uplink Radio Blocks starting with the next uplink Radio Block.

#### 6.6.4.1.1.2         Multiplexing of GPRS and EGPRS MSs

The GPRS and EGPRS MSs can be multiplexed dynamically on the same PDCH by utilising the USF. When uplink resources are allocated to a GPRS mobile, the network must use GMSK, i.e. CS-1 to CS-4 or MCS-1 to MCS-4, and the USF must point to the sequence of four uplink Radio Blocks starting with the next uplink Radio Block.

The dynamic allocation using USF granularity requires that a GPRS MS can read the USF in an EGPRS GMSK block. This is enabled by setting the stealing bits in the EGPRS GMSK blocks to indicate CS-4. The coding and interleaving of the USF is done as defined for CS-4. This leads to:

1.  A standard GPRS MS will be able to detect the USF in EGPRS GMSK blocks. The risk that the rest of the block will be misinterpreted as valid information is assumed to be low.

2.  An EGPRS MS can not differentiate CS-4 blocks and EGPRS GMSK blocks by only looking at the stealing bits. This is however not needed for USF detection, since the USF is signalled in the same way. Further, assuming that the EGPRS MS knows if it is in EGPRS or standard GPRS mode, it will only have to try to decode the remainder of the GMSK blocks in one way in order to determine if they were aimed for it.

A mobile station in EGPRS TBF mode shall be able to detect the USF that assigns the uplink to that mobile station. The network may use either GMSK modulation or 8-PSK modulation, i.e. CS-1 to CS-4, MCS-1 to MCS-4 or MCS-5 to MCS-9 in those blocks.

For mobile station synchronization reasons, if GPRS MSs are multiplexed on the PDCH, at least one downlink radio block every 360ms shall be transmitted to each MS with a coding scheme and a modulation that can be decoded by that MS.

### 6.6.4.1.2        Fixed Allocation

Fixed allocation where the Uplink part of the PDCH is reserved only for one MS during a certain period of time can be used to multiplex GPRS and EGPRS MSs on the same PDCH on the Uplink.

For MS synchronization reasons, if standard GPRS MSs are multiplexed on the PDCH, at least one Radio Block every 360ms on the Downlink must use GMSK (i.e. standard GPRS or MCS-1 to MCS-4).

### 6.6.4.1.3        Exclusive Allocation

Exclusive allocation is used to reserve the uplink part of the PDCH for only one MS during the life of the Temporary Block Flow. In exclusive allocation, all the uplink blocks of the uplink part of the PDCH are available to the MS for transmission.

### 6.6.4.2        Temporary Block Flow

A Temporary Block Flow (TBF) is a physical connection used by the two RR entities to support the unidirectional transfer of LLC PDUs on packet data physical channels. The TBF is allocated radio resource on one or more PDCHs and comprise a number of RLC/MAC blocks carrying one or more LLC PDUs. A TBF is temporary and is maintained only for the duration of the data transfer.

### 6.6.4.3        Temporary Flow Identity

Each TBF is assigned a Temporary Flow Identity (TFI) by the network. The assigned TFI is unique among concurrent TBFs in each directions and is used instead of the  MS identity in the RLC/MAC layer. The same TFI value may be used concurrently for TBFs in opposite directions. The TFI is assigned in a resource assignment message that precedes the transfer of LLC frames belonging to one TBF to/from the MS. The same TFI is included in every RLC header belonging to a particular TBF as well as in the control messages associated to the LLC frame transfer (e.g. acknowledgements) in order to address the peer RLC entities.

### 6.6.4.4        Medium Access modes

Four medium access modes are supported:

-   Dynamic allocation,

-   Extended Dynamic allocation,

-   Fixed allocation, and

-   Exclusive allocation.

The Dynamic allocation medium access mode or Fixed allocation medium access mode shall be supported by all networks that support GPRS. The support of Extended Dynamic allocation and Exclusive allocation medium access modes is optional.

The Dynamic allocation and Fixed allocation modes shall be supported in all mobile stations. Exclusive allocation shall be supported in all mobile stations supporting DTM.

### 6.6.4.5        Acknowledged mode for RLC/MAC operation

### 6.6.4.5.1        GPRS

The transfer of RLC Data Blocks in the acknowledged RLC/MAC mode is controlled by a selective ARQ mechanism coupled with the numbering of the RLC Data Blocks within one Temporary Block Flow. The sending side (the MS or the network) transmits blocks within a window and the receiving side sends Packet Uplink Ack/Nack or Packet Downlink Ack/Nack message when needed. Every such message acknowledges all correctly received RLC Data Blocks

up to an indicated block sequence number (BSN), thus "moving" the beginning of the sending window on the sending side. Additionally, the bitmap that starts at the same RLC Data Block is used to selectively request erroneously received RLC Data Blocks for retransmission. The sending side then retransmits the erroneous RLC Data Blocks, eventually resulting in further sliding the sending window.

The Packet Ack/Nack message does not include any change in the current assignment (and thus does not have to be acknowledged when sent on downlink). A missing Packet Ack/Nack is not critical and a new one can be issued whenever. In Packet Downlink Ack/Nack message, the MS may optionally initiate an uplink TBF. In Packet Uplink Ack/Nack message , the network can assign uplink resources for mobile station using a fixed allocation.

When receiving uplink data from a MS the network shall, based on erroneous blocks received from MS, allocate additional resources for retransmission.

The acknowledgement procedure of the LLC layer is not combined with the acknowledgement procedure on the underlying RLC/MAC layer.

## 6.6.4.5.2          EGPRS

The transfer of RLC Data Blocks in the acknowledged RLC/MAC mode can be controlled by a selective type I ARQ mechanism, or by type II hybrid ARQ (incremental redundancy: IR) mechanism, coupled with the numbering of the RLC Data Blocks within one Temporary Block Flow. The sending side (the MS or the network) transmits blocks within a window and the receiving side sends Packet Uplink Ack/Nack or Packet Downlink Ack/Nack message when needed.

According to the link quality, an initial MCS is selected for an RLC block. For the retransmissions, the same or another MCS from the same family of MCSs can be selected. E.g. if MCS-7 is selected for the first transmission of an RLC block, any MCS of the family B can be used for the retransmissions. The selection of MCS is controlled by the network.

In the EGPRS type II Hybrid ARQ scheme, the information is first sent with one of the initial code rates (i.e., the rate 1/3 encoded data is punctured with the puncturing scheme (PS) 1 of the selected MCS). If the RLC Data Block is received in error, additional coded bits (i.e., the output of the rate 1/3 encoded data which is punctured with PS 2 of the prevailing MCS) are sent and decoded together with the already received codewords until decoding succeeds. If all the codewords (different punctured versions of the encoded data block) have been sent, the first codeword (which is punctured with PS 1) is sent. Alternatively, it is possible to use incremental redundancy modes called MCS-5-7 and MCS-6-9, in which the initial transmissions are sent with either MCS-5 or MCS-6 (respectively) and the retransmissions are sent with MCS-7 or MCS-9 (respectively). Header part is robustly coded so that the receiver is able to determine the block identities for all transmissions, even if the payload cannot be decoded.

In the EGPRS type I ARQ, the operation is similar to the one of the EGPRS type II hybrid ARQ, except that the decoding of an RLC Data Block is solely based on the prevailing transmission (i.e., erroneous blocks are not stored).

Type II hybrid ARQ is mandatory in EGPRS MS receivers and the associated performance requirements are specified in GSM 05.05 [14].

## 6.6.4.6          Unacknowledged mode for RLC/MAC operation

The transfer of RLC Data Blocks in the unacknowledged RLC/MAC mode is controlled by the numbering of the RLC Data Blocks within one Temporary Block Flow and does not include any retransmissions. The receiving side extracts user data from the received RLC Data Blocks and attempts to preserve the user information length by replacing missing RLC Data Blocks by dummy information bits.

The same mechanism and message format for sending temporary acknowledgement messages is used as for acknowledged mode in order to convey the necessary control signalling (e.g. monitoring of channel quality for downlink channel or timing advance correction for uplink transfers). The fields for denoting the erroneous RLC blocks may be used as an additional measure for channel quality (i.e. parameter for link adaptation). The sending side (the MS or the network) transmits a number of radio blocks and then polls the receiving side to send an acknowledgement message. The Packet Uplink Ack/Nack and Packet Downlink Ack/Nack message does not include any change in the current assignment. A missing acknowledgement message is not critical and a new one can be obtained whenever. In Packet Downlink Ack/Nack message, the MS may optionally initiate an uplink TBF. In Packet Uplink Ack/Nack message , the network can assign uplink resources for mobile station using a fixed allocation.

### 6.6.4.7 Mobile Originated Packet Transfer

#### 6.6.4.7.1 Uplink Access

##### 6.6.4.7.1.1 On the (P)RACH

This sub-clause applies to all mobile stations in packet idle mode. It also applies to class A mobiles stations in dedicated mode if they are not DTM capable.
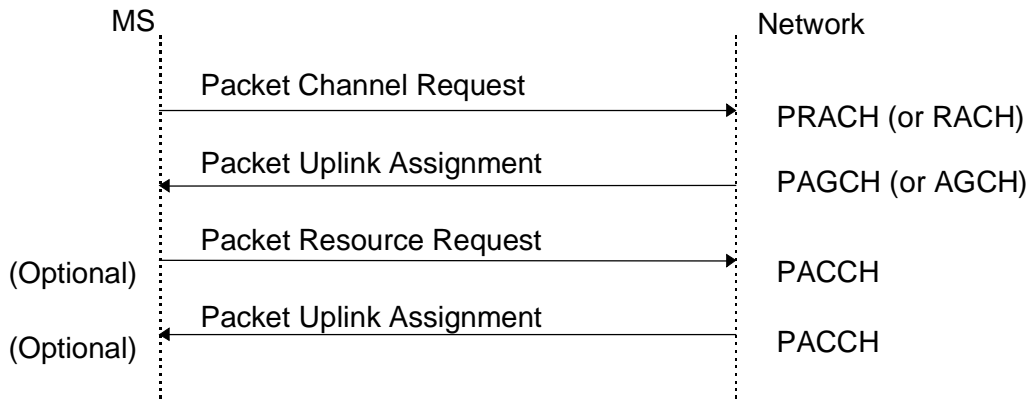


**Figure 23: Access and allocation for the one or two phase packet access, uplink packet transfer**

An MS initiates a packet transfer by making a Packet Channel Request on PRACH or RACH. The network responds on PAGCH or AGCH respectively. It is possible to use one or two phase packet access method (see Figure 23).

In the one phase access, the Packet Channel Request is responded by the network with the Packet Uplink Assignment reserving the resources on PDCH(s) for uplink transfer of a number of Radio blocks. The reservation is done accordingly to the information about the requested resources that is comprised in the Packet Channel Request. On RACH, there is only two cause values available for denoting GPRS, which can be used to request limited resources or two phase access. On PRACH, the Packet Channel Request may contain more adequate information about the requested resources and, consequently, uplink resources on one or several PDCHs can be assigned by using the Packet Uplink Assignment message.

In the two phase access, the Packet Channel Request is responded with the Packet Uplink Assignment which reserves the uplink resources for transmitting the Packet Resource Request. A two phase access can be initiated by the network or a mobile station. The network can order the MS to send Packet Resource Request message by setting parameter in Packet Uplink Assignment message. Mobile station can require two phase access in Packet Channel Request message. In this case, the network may order MS to send Packet Resource Request or continue with a one phase access procedure.

The Packet Resource Request message carries the complete description of the requested resources for the uplink transfer. The MS can indicate the medium access method, it prefers to be used during the TBF. The network responds with the Packet Uplink Assignment reserving resources for the uplink transfer and defining the actual parameters for data transfer (e.g. medium access mode).

If there is no response to the Packet Channel Request within predefined time period, the MS makes a retry after a random backoff time.

On PRACH there is used a 2-step approach including a long-term and a short-term estimation of the persistence (see Figure 24). The optimal persistence of the mobile stations is calculated at the network side.
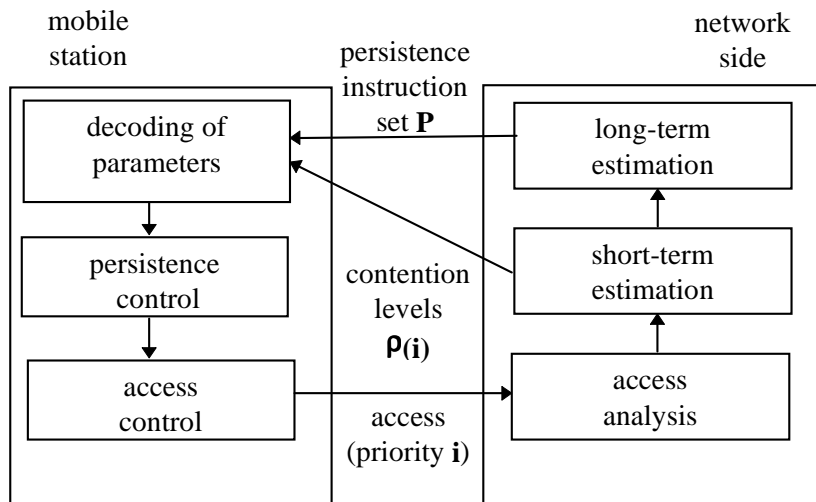
**Figure 24: Basic principle of random access traffic control**

The actual persistence values depend on:

-   the priority i of the packet to be transmitted;

-   the amount of traffic within higher priority classes;

-   the amount of traffic within the own priority class.

Optionally, the existing backoff algorithm on RACH can be used on PRACH.

On RACH, the existing backoff algorithm shall be used.

Occasionally, more Packet Channel Requests can be received than can be served. To handle this, a Packet Queuing Notification is transmitted to the sender of the Packet Channel Request. The notification includes information that the Packet Channel Request message is correctly received and Packet Uplink Assignment may be transmitted later. If the Timing Advance information becomes inaccurate for an MS, the network can send Packet Polling Request to trigger the MS to send four random access bursts. This can be used to estimate the new Timing Advance before issuing the Packet Uplink Assignment.

6.6.4.7.1.2          On the main DCCH

This clause only applies to mobile stations in dedicated mode that support DTM. This procedure moves the mobile station from dedicated mode to dual transfer mode.



**Figure 24b: Access and allocation for the packet request procedure, uplink packet transfer**

A DTM mobile station initiates a packet transfer while in dedicated mode by sending a DTM Request message on the main DCCH.

The DTM Request message carries a description of the requested resources for the uplink transfer. The DTM Request message is responded by the network with one of the following  DTM assignment messages:

- DTM Assignment Command: when the network allocates a TBF and reallocation of the resource of the RR connection is needed.

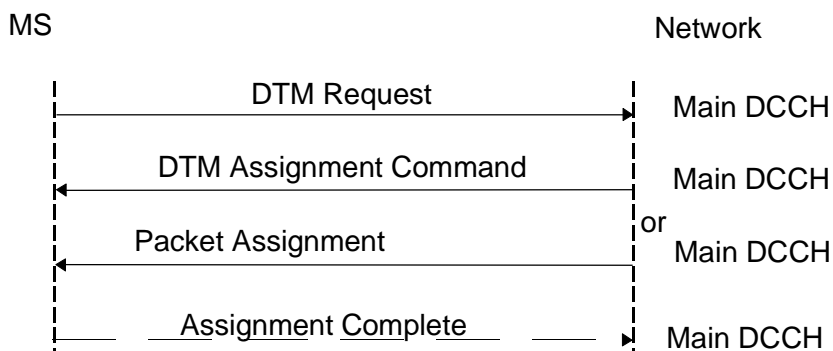- Packet Assignment: when the network allocates a TBF and no reallocation of the resource of the RR connection is needed.

If there is reallocation of the resource of the RR connection (through a DTM Assignment Command message), the MS sends an Assignment Complete message on the new main DCCH after it is established.

### 6.6.4.7.2 Dynamic/Extended Dynamic allocation

#### 6.6.4.7.2.1 Uplink Packet Transfer

The Packet Uplink Assignment message includes the list of PDCHs and the corresponding USF value per PDCH. A unique TFI is allocated and is thereafter included in each RLC Data and Control Block related to that Temporary Block Flow. The MS monitors the USFs on the allocated PDCHs and transmits Radio blocks on those which currently bear the USF value reserved for the usage of the MS.

If the resource assigned by the network in the case of extended dynamic allocation does not allow the multislot MS (see GSM 05.02 [11], annex B) to monitor the USF on all the assigned PDCHs, the following rules shall apply:

- Whenever the MS receives its USF on one downlink PDCH (e.g. on timeslot 0 while timeslots 0, 2 and 3 were assigned), it shall consider the corresponding uplink block and all subsequent ones from the list of assigned PDCHs as allocated (e.g. on 0, 2 and 3). Hence, if the network allocates a block to this MS on an assigned PDCH, it shall also allocate blocks to this MS on all subsequent PDCHs in the list. For each allocated block, the network shall set the USF to the value reserved for the usage of that MS. These rules apply on a block period basis.

- During block periods where it is transmitting, the MS shall monitor the USF on each PDCH in the list of assigned PDCHs, up to and including the first PDCH currently used for transmission. This rule applies on a block period basis. For example, if timeslots 0, 2 and 3 have been assigned and blocks are currently allocated on timeslots 2 and 3, then during this block period the MS monitors USF on timeslots 0 and 2. If the reserved value of USF is found on timeslot 0, then the next allocated blocks shall be on timeslots 0, 2 and 3. If the reserved value of USF is found on timeslot 2, then the next allocated blocks shall be on timeslots 2 and 3. And so on for the subsequent block periods. Because each Radio Block includes an identifier (TFI), all received Radio blocks are correctly associated with a particular LLC frame and a particular MS, thus making the protocol highly robust. By altering the state of USF, different PDCHs can be "opened" and "closed" dynamically for certain MSs thus providing a flexible reservation mechanism. Additionally, packets with higher priority and pending control messages can temporarily interrupt a data transmission from one MS.

The channel reservation algorithm can also be implemented on assignment basis. This allows individual MSs to transmit a predetermined amount of time without interruptions.

The MS may be allowed to use the uplink resources as long as there is queued data on the RLC/MAC layer to be sent from the MS. It can comprise a number of LLC frames. In that sense the radio resources are assigned on the initially "unlimited" time basis. Alternatively, the uplink assignment for each assignment may be limited to a number of radio blocks (e.g. in order to offer more fair access to the medium at higher loads).

The selective ARQ operation for the acknowledged RLC/MAC mode is described in Subclause 6.6.4.5. The unacknowledged RLC/MAC mode operation is described in Subclause 6.6.4.6.

Figure 25 shows an example of message sequence for the (multislot) uplink data transfer with one resource reallocation and possible RLC Data Block re-transmissions.
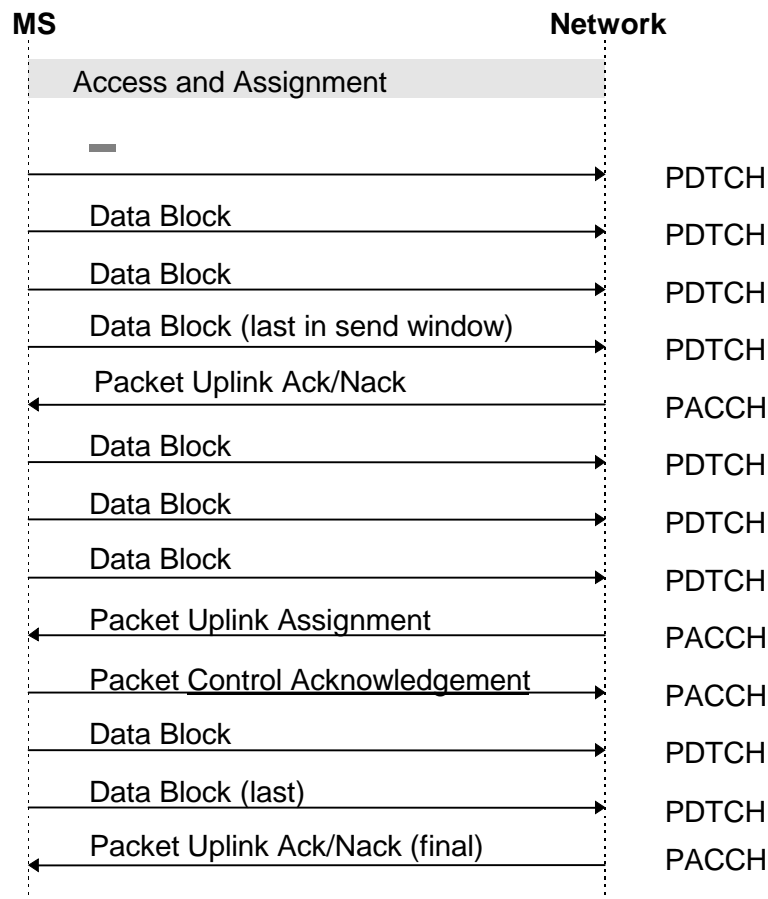
**MS**                      **Network**

| | |
|---|---|
| Access and Assignment | |
| | PDTCH |
| Data Block | PDTCH |
| Data Block | PDTCH |
| Data Block (last in send window) | PDTCH |
| Packet Uplink Ack/Nack | PACCH |
| Data Block | PDTCH |
| Data Block | PDTCH |
| Data Block | PDTCH |
| Packet Uplink Assignment | PACCH |
| Packet Control Acknowledgement | PACCH |
| Data Block | PDTCH |
| Data Block (last) | PDTCH |
| Packet Uplink Ack/Nack (final) | PACCH |

**Figure 25: An example of dynamic allocation uplink data transfer**

6.6.4.7.2.2          Release of the Resources

The release of the resources is normally initiated from the MS by counting down the last couple of blocks.

For the normal release of resources for RLC connection carrying a mobile originated packet transfer, the mechanism based on acknowledged final Packet Uplink Ack/Nack combined with timers is used.

After the MS has sent its last RLC Data Block (indicated by the countdown field), the acknowledgement is expected from the network side. By sending the last block, the MS may no longer use the same assignment unless a negative acknowledgement arrives. It also means that the network side may reallocate the same USF(s) to some other user as soon as all the RLC Data Blocks belonging to that Temporary Block Flow are correctly received; that regardless of the possible later errors in the acknowledgements.

The next step, in the case of all RLC Data Blocks being correctly received, is that the network sends Packet Uplink Ack/Nack which is to be immediately acknowledged by the MS in the reserved uplink block period. It must be possible for the network not to use the mechanism of acknowledgement for Packet Ack/Nack in which case the release of the resources procedure relies only on timers. The TFI can be reused for another assignment either upon the reception of the acknowledgement for Packet Ack/Nack or after expiry of the guard timer.

Further, the premature release or change of assignment for one MS may be initiated:

- by the network with an explicit message

- in dual transfer mode, by the release of the RR connection (i.e. handover, assignment and channel release procedures) or

- by the establishment of an RR connection.

In the case of release, the MS is ordered to interrupt the Temporary Block Flow. The MS shall then reorganise the uplink buffer and issue a new Packet Channel Request to continue the uplink transfer with the RLC Data Blocks containing untransferred (i.e. on the RLC/MAC layer unacknowledged) LLC frames.

A change in assignment may also be initiated by the network, in which case the Packet Uplink Assignment, Packet Timeslot Reconfigure, DTM Assignment Command or Packet Assignment message is issued.

### 6.6.4.7.3 Fixed Allocation

Fixed allocation uses the Packet Uplink Assignment message to communicate a detailed fixed uplink resource allocation to the MS. The fixed allocation consists of a start frame, slot assignment, and block assignment bitmap representing the assigned blocks per timeslot. The MS waits until the start frame indicated and then transmits radio blocks on those blocks indicated in the block assignment bitmap. The fixed allocation does not include the USF and the MS is free to transmit on the uplink without monitoring the downlink for the USF. Unused USF value is used to prevent other mobiles to transmit. If the current allocation is not sufficient, the MS may request additional resources in one of the assigned uplink blocks. A unique TFI is allocated and is thereafter included in each RLC data and control block related to that Temporary Block Flow. Because each Radio Block includes an identifier (TFI), all received Radio blocks are correctly associated with a particular LLC frame and a particular MS.

The number of blocks an MS requests in the initial and subsequent allocation request messages shall only account for the number of data and control blocks it intends to send. The MS shall not request additional blocks for the retransmission of erroneous blocks. The network can repeat the allocation of radio resources by setting the parameter in the Packet Uplink Assignment or the Packet Uplink Ack/Nack message.

The selective ARQ operation for the acknowledged RLC/MAC mode is described in Subclause 6.6.4.5. The unacknowledged RLC/MAC mode operation is described in Subclause 6.6.4.6.

Figure 25 shows an example of message sequence for the (multislot) uplink data transfer with one resource reallocation and possible RLC Data Block re-transmissions.

### 6.6.4.7.4 Exclusive Allocation

Exclusive allocation uses the Packet Uplink Assignment, Packet Timeslot Reconfigure, DTM Assignment Command or Packet Assignment messages to communicate a exclusive uplink resource allocation to the MS. The exclusive allocation consists of a start frame and a slot assignment. The MS waits until the start frame indicated and then starts transmitting the radio blocks. The exclusive allocation includes neither the USF nor the block assignment bitmap. The MS is free to transmit on the uplink without monitoring the downlink for the USF. Unused USF values may be used to prevent other mobiles from transmitting. A unique TFI is allocated and is thereafter included in each RLC data and downlink control block related to that Temporary Block Flow. Because each Radio Block includes an identifier (TFI or TLLI), all received Radio blocks are correctly associated with a particular LLC frame and a particular MS.

For a close ended TBF, the number of blocks an MS requests in the initial and subsequent allocation request messages shall only account for the number of data blocks it intends to send. The MS shall not request additional blocks for the retransmission of erroneous blocks.

The selective ARQ operation for the acknowledged RLC/MAC mode is described in Subclause 6.6.4.5. The unacknowledged RLC/MAC mode operation is described in Subclause 6.6.4.6.

Figure 25 shows an example of message sequence for the (multislot) uplink data transfer with one resource reallocation and possible RLC Data Block re-transmissions.

### 6.6.4.7.5 Contention Resolution

Contention resolution is an important part of RLC/MAC protocol operation, especially because one channel allocation can be used to transfer a number of LLC frames. Contention resolution applies for both dynamic and fixed allocation medium access modes. Contention resolution does not apply to the packet request procedures while in dedicated mode for an MS supporting DTM.

There are two basic access possibilities, one phase and two phase access as defined in Subclause 6.6.4.7.1.

The two phase access is inherently immune for possibility that two MSs can perceive the same channel allocation as their own. Namely the second access phase, the Packet Resource Request, uniquely identifies the MS by its TLLI. The same TLLI is included in the Packet Uplink Assignment/Packet Downlink Assignment and no mistake is possible.

The one phase access is somewhat insecure and an efficient contention resolution mechanism has to be introduced.

The first part of the solution is the identification of the MS. The identification of transmitting MS on the RLC/MAC level is necessary not only for contention resolution but also to be able to establish RLC protocol entity for that Temporary Block Flow on the network side. Additionally, the TLLI is necessary to be able to match simultaneous uplink and downlink packet transfers by taking into consideration multislot capability of that MS.

In order to uniquely identify the MS when sending on uplink, the RLC Header for all the RLC Data Blocks on uplink is extended to include the TLLI until the contention resolution is completed on the MS side.

The second part of the solution is the notification from the network side about who owns the allocation. That is solved by the inclusion of the TLLI in the Packet Uplink Ack/Nack/Packet Downlink Ack/Nack. This message shall be sent in an early stage, even before the receive window for RLC/MAC protocol operation is full. By doing so, the contention is resolved after the first occurrence of Packet Ack/Nack. The possibility of RLC Data Blocks being captured from "wrong" MS, thus destroying the LLC frame, shall be covered for by retransmissions on the LLC layer.

## 6.6.4.8 Mobile Terminated Packet Transfer

### 6.6.4.8.1 Packet Paging

The network initiates a packet transfer to an MS that is in the Standby state by sending one or more packet paging request messages on the downlink PPCH or PCH. The MS responds to one packet paging request message by initiating a mobile originated packet transfer, as described in section 6.6.4.7. This mobile originated packet transfer allows the MS to send a packet paging response to the network. The packet paging response is one or more RLC/MAC data blocks containing an arbitrary LLC frame. The message sequence described in Figure 26 below is conveyed either on PCCCH or on CCCH. After the packet paging response is sent by the MS and received by the network, the mobility management state of the MS is Ready.

The network can then assign some radio resources to the MS and perform the downlink data transfer as described in section 6.6.4.8.2.
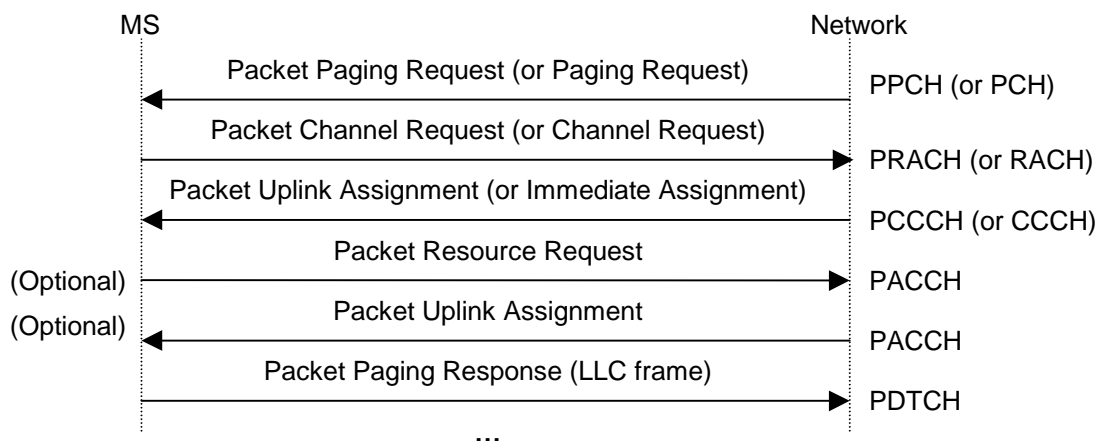


**Figure 26: Paging message sequence for Paging, downlink packet transfer**

### 6.6.4.8.2 Downlink Packet Transfer

The transmission of a packet to an MS in the Ready state is initiated by the network using a packet downlink assignment message. In case there is an uplink packet transfer in progress, the packet downlink assignment message is transmitted on PACCH. Else, in case there is PCCCH allocated in the cell, the Packet Downlink Assignment message is transmitted on PCCCH. Else, the Immediate Assignment message is transmitted on CCCH. Else, for a DTM MS in dedicated mode, the assignment message is transmitted on the main DCCH. The packet downlink assignment message includes the list of PDCH(s) that will be used for downlink transfer. The Timing Advance and Power Control information is also included, if available. Otherwise, the MS may be requested to respond with a Packet Control Acknowledgement (see also Subclause 6.5.7 on timing advance procedures). The MS multislot capability needs to be considered.

The network sends the RLC/MAC blocks belonging to one Temporary Block Flow on downlink on the assigned downlink channels.

Multiplexing the RLC/MAC blocks destined for different MSs on the same PDCH downlink is enabled with an identifier, e.g. TFI, included in each RLC/MAC block. The interruption of data transmission to one MS is possible.

The acknowledged (i.e. selective ARQ operation) and unacknowledged RLC/MAC mode operation is described in Subclauses 6.6.4.5 and 6.6.4.6. The sending of the Packet Downlink Ack/Nack message is obtained by the occasional network initiated polling of the MS. The MS sends the Packet Downlink Ack/Nack message in a reserved radio block which is allocated together with polling. Unassigned USF value is used in the downlink radio block which corresponds to the reserved uplink radio blocks. Further, if the MS wants to send some additional signalling or uplink data, it may be indicated in the Packet Downlink Ack/Nack message.

Figure 27 shows an example of message sequence for (multislot) downlink data transfer with one resource reallocation and possible RLC Data Block re-transmissions.
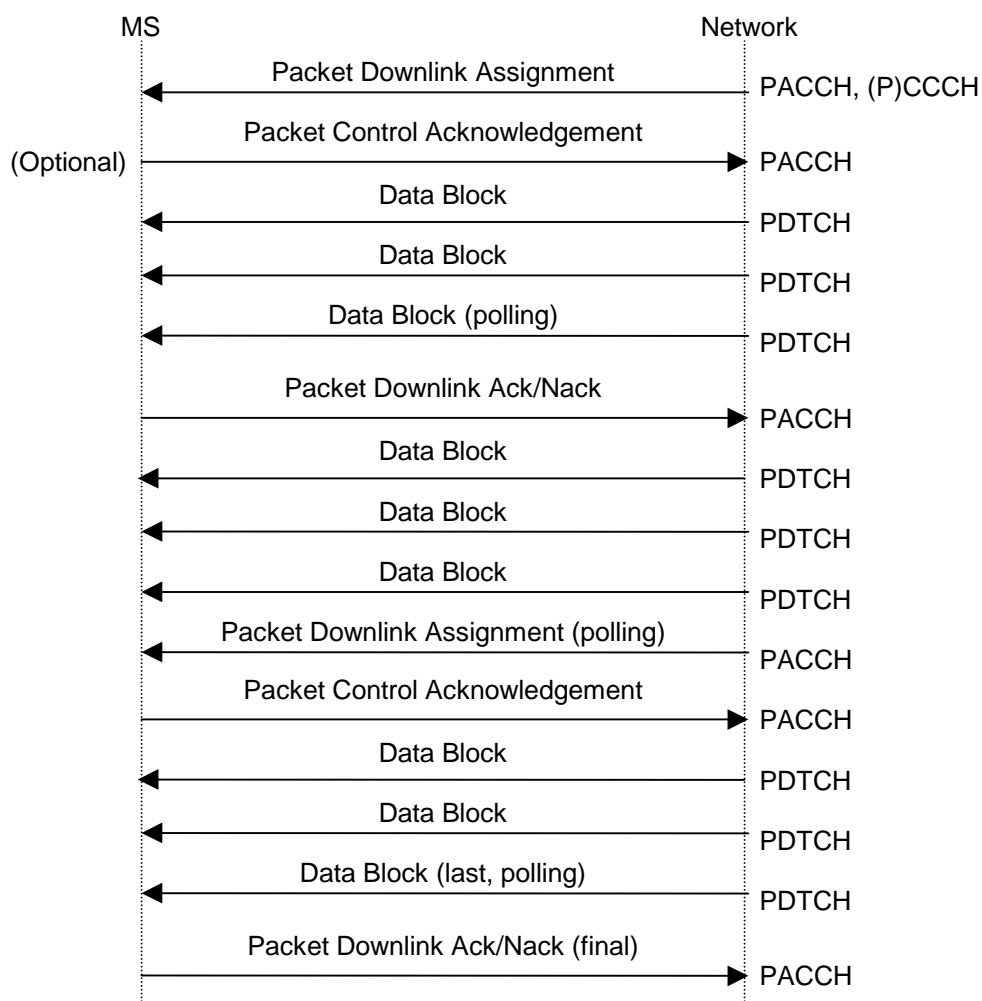


**Figure 27: An example of downlink data transfer**

### 6.6.4.8.3          Release of the Resources

The release of the resources is initiated by the network by terminating the downlink transfer and polling the MS for a final Packet Downlink Ack/Nack message.

A mobile station in dual transfer mode shall abandon the packet resources when the RR connection is released.

It is possible for the network to change the current downlink assignment by using the Packet Downlink Assignment or Packet Timeslot Reconfigure message, which then has to be acknowledged by the MS in a reserved radio block on the uplink.

The handling of TFI is steered with the same timer that runs on both the MS and the network side after the last RLC Data Block is sent to the MS. When it expires, the current assignment becomes invalid for the MS and TFI can be reused by the network. Further, upon the reception of the final Packet Downlink Ack/Nack from the MS, another timer is started on network side. When it expires, the current assignment becomes invalid for the MS and TFI can be reused by the network.

### 6.6.4.8.4 Packet Paging Notification

The network initiates a packet transfer to a DTM MS that is in the Standby state and in dedicated mode by sending one or more packet paging notification messages on the downlink main DCCH. The MS responds to one packet paging notification message by initiating a mobile originated packet request, as described in section 6.6.4.7. This mobile originated packet transfer allows the MS to send a packet paging response to the network. The packet paging response is one or more RLC/MAC data blocks containing an arbitrary LLC frame. After the packet paging response is sent by the MS and received by the network, the mobility management state of the MS is Ready.

The network can then assign some radio resources to the MS and perform the downlink data transfer as described in section 6.6.4.8.2.

### 6.6.4.9 Simultaneous Uplink and Downlink Packet Transfer

During the ongoing uplink Temporary Block Flow, the MS continuously monitors one downlink PDCH for possible occurrences of Packet Downlink Assignment or Packet Timeslot Reconfigure messages on PACCH (see Figure 25). The MS is therefore reachable for downlink packet transfers that can then be conveyed simultaneously on the PDCH(s) that respect the MS multislot capability.

If the MS wants to send packets to the network during the ongoing downlink Temporary Block Flow, it can be indicated in the acknowledgement that is sent from the MS. By doing so, no explicit Packet Channel Requests have to be sent to the network. Further, the network already has the knowledge of which PDCH(s) that particular MS is currently using so that the uplink resources can be assigned on the PDCH(s) that respect the MS multislot capability. This method may introduce an extra delay when initiating the uplink packet transfer but only for the first LLC frame in a sequence.

# 6.7 Abnormal cases in GPRS MS Ready State

The RLC/MAC error causes and procedures to handle these can be found in GSM 04.08 [6] , 04.60 [7] and 05.08 [15].

# 6.8 PTM-M Data Transfer

> NOTE: The stage 3 specification for PTM-M data transfer is left for phase 2 of GPRS specification.

PTM-M data, in the form of individual LLC frames, is mapped into RLC/MAC-PTM_DATA primitive and distributed from SGSN to the BSS representing the cells that are defined by a geographical area parameter. To the cells concerned, the BSS for each PTM-M LLC frame:

- Optionally, sends a "PTM-M new message" indicator on all individual paging channels on PCCCH if allocated, otherwise on CCCH. The indication refers to a PTM-M notification channel PNCH on PCCCH or NCH on CCCH, where a notification for the new PTM-M message can be received.

  If the indicator option is not supported, or if an MS can not receive the indicator when expected, e.g. because the corresponding block in the multiframe structure is used for other purposes than paging, the MS must read the notification channel.

- Sends a PTM-M notification on PNCH or NCH. The notification has the form of a Packet Resource Assignment for the PTM-M LLC frame. The notification includes a group identity IMGI, a unique LLC frame identifier (in the form of an N-PDU number together with a segment offset, see GSM 04.65) and an allocation of a TFI to be used in all RLC blocks of the LLC frame.

- Transmits the PTM-M LLC frame on the assigned downlink resources.

Transfer of PTM-M data is carried out without any ARQ on the RLC/MAC  and LLC layers. Instead, each LLC frame is retransmitted a specified number of times. For each retransmission, the above procedure is performed.  The PTM-M notification (resource assignment)  includes the unique LLC frame identifier as in the first transmission but a new allocation of TFI.

An MS accumulates correctly received RLC blocks from each transmission to assemble an LLC frame.

The dimensioning of PNCH shall be scaleable depending on capacity requirements.

An NCH may, if capacity allows, be used as a shared notification channel for PTM-M and Advanced Speech Call Items (ASCI).

An MS only interested in PTM-M needs to listen only to PNCH/NCH.

# Annex A (informative):
# Bibliography

1) ITU-T I.130, Method for the Characterization of Telecommunication Services Supported by an ISDN

2) ITU-T Q.65, Stage 2 of the Method for Characterization of the Services Supported by an ISDN

3) DIS 8886, OSI Data Link Service Definition

4) DIS 10022, OSI Physical Service Definition

5) ISO 10039, Medium Access Control Service Definition

6) ISO 4335, HDLC Procedures

7) ISO 7478, Multilink Procedures

8) ISO 7498, OSI Basic Reference Model and Layer Service Conventions

# Annex B (informative):
# Change history

| SPEC | SMG# | CR | PHA | VERS | NEW_VER | SUBJECT |
|------|------|-----|-----|------|---------|---------|
| 03.64 | s22 | NEW | 2+ | 2.1.1 | 5.0.0 | GSM 03.64 GPRS Stage 2 Radio |
| 03.64 | s23 | A022 | R97 | 5.0.0 | 5.1.0 | Unacknowledged mode of RLC/MAC operation |
| 03.64 | s23 | A023 | R97 | 5.0.0 | 5.1.0 | Improved RLC Service Primitives |
| 03.64 | s23 | A024 | R97 | 5.0.0 | 5.1.0 | Enhancements to dynamic allocation |
| 03.64 | s23 | A025 | R97 | 5.0.0 | 5.1.0 | Clarifications to DRX |
| 03.64 | s23 | A026 | R97 | 5.0.0 | 5.1.0 | Optimisation for network control cell reselection |
| 03.64 | s23 | A027 | R97 | 5.0.0 | 5.1.0 | Abnormal Cases in GPRS MS Ready State: Leaky Bucket Procedure |
| 03.64 | s23 | A029 | R97 | 5.0.0 | 5.1.0 | Multiframe structure (details) (revision of SMG2 GPRS 301/97) |
| 03.64 | s23 | A030 | R97 | 5.0.0 | 5.1.0 | Abnormal Cases in GPRS MS Ready State |
| 03.64 | s23 | A031 | R97 | 5.0.0 | 5.1.0 | Cell Re-Selection in GPRS |
| 03.64 | s23 | A032 | R97 | 5.0.0 | 5.1.0 | Definition of PACCH |
| 03.64 | s23 | A033 | R97 | 5.0.0 | 5.1.0 | Clarifications on Timing advance procedure |
| 03.64 | s23 | A035 | R97 | 5.0.0 | 5.1.0 | Bit order for USF coding in GPRS |
| 03.64 | s23 | A036 | R97 | 5.0.0 | 5.1.0 | PTM-M |
| 03.64 | s23 | A037 | R97 | 5.0.0 | 5.1.0 | Contention resolution |
| 03.64 | s23 | A039 | R97 | 5.0.0 | 5.1.0 | Deleting parameter XHYST |
| 03.64 | s24 | A031 | R97 | 5.1.0 | 5.2.0 | Clarification on the use of hysteresis for cell re-selection |
| 03.64 | s25 | A043 | R97 | 5.2.0 | 6.0.0 | Clarification of the use of TAI |
| 03.64 | s25 | A049 | R97 | 5.2.0 | 6.0.0 | USF granularity for dynamic allocation |
| 03.64 | s26 | | R97 | 6.0.0 | 6.0.1 | Editorial changes for Publication |
| 03.64 | s27 | A050 | R97 | 6.0.1 | 6.1.0 | Changes on all chapters to align the spec. with other GPRS specifications. |
| 03.64 | s27 | A044 | R97 | 6.0.1 | 6.1.0 | Clarification on PACCH allocation for fixed assignment |
| 03.64 | s28 | A052 | R97 | 6.1.0 | 6.2.0 | PBCCH scheduling and editorial corrections |
| 03.64 | s28 | A051 | R97 | 6.1.0 | 6.2.0 | Interference measurements on network control |
| 03.64 | s28 | A053 | R97 | 6.1.0 | 6.2.0 | 51-multiframe PBCCH |
| 03.64 | s29 | A055 | R97 | 6.2.0 | 6.3.0 | Miscellaneous corrections |
| 03.64 | s29 | A056 | R97 | 6.2.0 | 6.3.0 | Clarification of polling response |
| 03.64 | s29 | A057 | R97 | 6.2.0 | 6.3.0 | Correction to 1 phase access contention resolution |
| 03.64 | s29 | A054 | R99 | 6.2.0 | 8.0.0 | Introduction of Enhanced GPRS (EGPRS) |
| 03.64 | s30 | A066 | 2 | 8.0.0 | 8.1.0 | MT packet transfer |
| 03.64 | s30 | A058 | | 8.0.0 | 8.1.0 | EGPRS fine tuning |
| 03.64 | s30 | A059 | 1 | 8.0.0 | 8.1.0 | EDGE Compact Cell Reselection |
| 03.64 | s30 | A060 | 1 | 8.0.0 | 8.1.0 | EDGE Compact logical channels |
| 03.64 | s30bis | A067 | | 8.1.0 | 8.2.0 | Correction for EGPRS |
| 03.64 | s30bis | A068 | | 8.1.0 | 8.2.0 | Introduction of Incremental Redundancy between different coding schemes in EGPRS |
| 03.64 | s31 | A069 | | 8.2.0 | 8.3.0 | GPRS & SMS-CB interworking |
| 03.64 | s31bis | A070 | R99 | 8.3.0 | 8.4.0 | 8-PSK support in UL |
| 03.64 | s32 | A072r2 | R99 | 8.4.0 | 8.5.0 | DTM: Stage 2 Description |
| | | | | | | September 2000 TSG-GERAN |
| 03.64 | G01 | A073 | 1 | 8.5.0 | 8.6.0 | CR 03.64-A073 rev 1 DTM (R99) |
| 03.64 | G01 | A074 | | 8.5.0 | 8.6.0 | CR 03.64-A074 DTM+EGPRS (R99) |
| 03.64 | G01 | A075 | 1 | 8.5.0 | 8.6.0 | CR 03.64-A075 rev 1 DTM: alignments (R99) |
| 03.64 | G03 | A077 | | 8.6.0 | 8.7.0 | Removal of Anonymous Access |
| 03.64 | G03 | A078 | 1 | 8.6.0 | 8.7.0 | Correction to the ordering of E/FBI bits |
| 03.64 | G04 | A080 | 1 | 8.7.0 | 8.8.0 | Applicability of GPRS requirements to EGPRS |
| 03.64 | G07 | A083 | | 8.8.0 | 8.9.0 | Coding rate of MCS-3 |
| 03.64 | G07 | A084 | | 8.8.0 | 8.9.0 | Clarification of EGPRS MS USF decoding |
| 03.64 | G08 | A085 | 1 | 8.9.0 | 8.10.0 | Network requirements for MS synchronisation |
| 03.64 | G14 | A086 | 1 | 8.10.0 | 8.11.0 | Padding for MCS-8 retransmissions |
| 03.64 | G19 | A088 | 2 | 8.11.0 | 8.12.0 | Change of DTM core capability |

# History

| Document history | | |
|---|---|---|
| V8.3.0 | April 2000 | Publication |
| V8.4.0 | June 2000 | Publication |
| V8.5.0 | August 2000 | Publication |
| V8.6.0 | September 2000 | Publication |
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